

Lecture 5

Subgame Perfect Nash Equilibrium
and Backward Induction in games of
imperfect information

Backward Induction in dynamic games of imperfect information

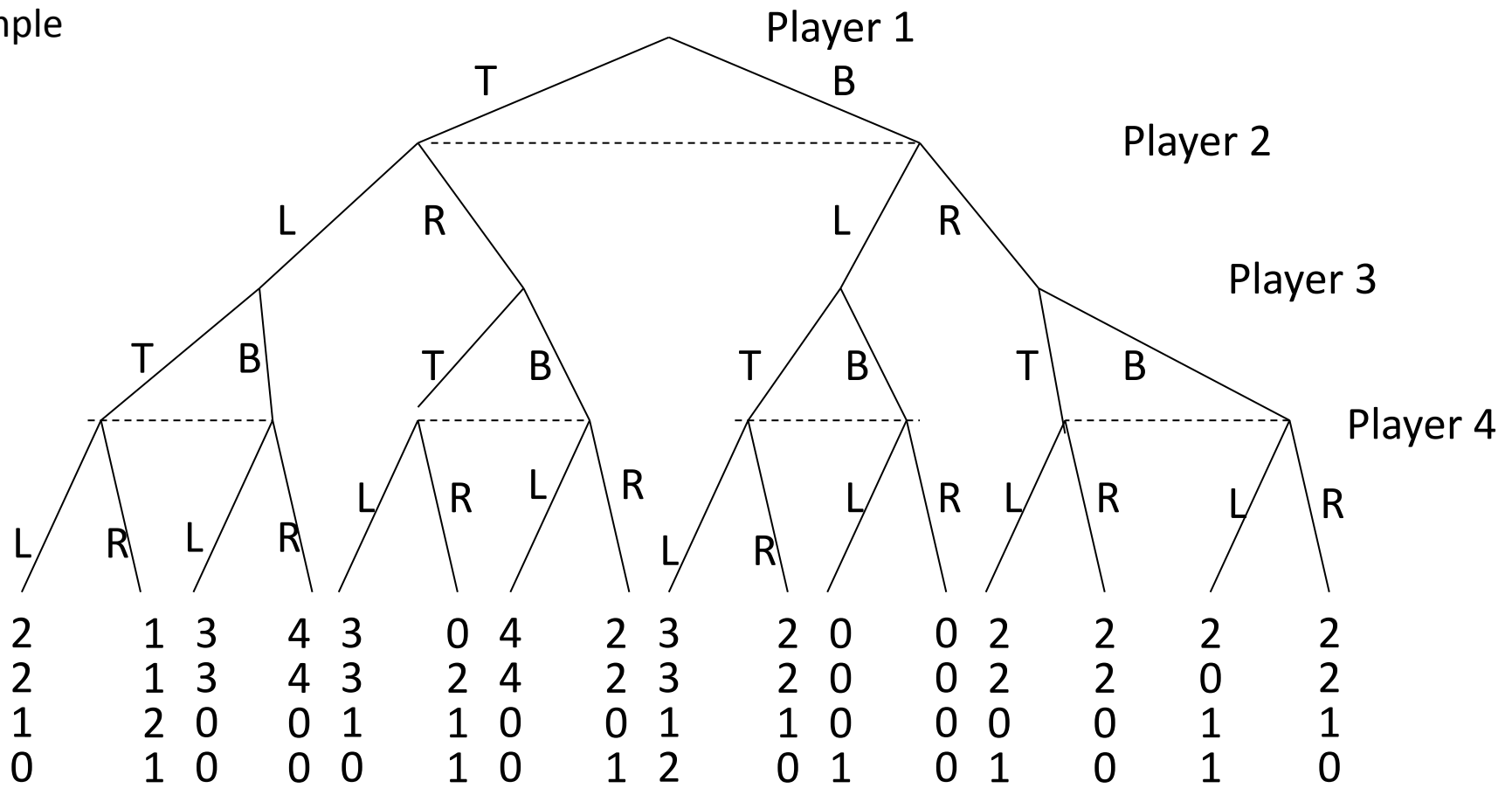
- We start at the end of the trees
- first find the Nash equilibrium (NE) of the last subgame
- then taking this NE as given, find the NE in the second last subgame
- continue working backwards

If in each subgame there is only one NE, this procedure leads to a **Unique Subgame Perfect Nash equilibrium**

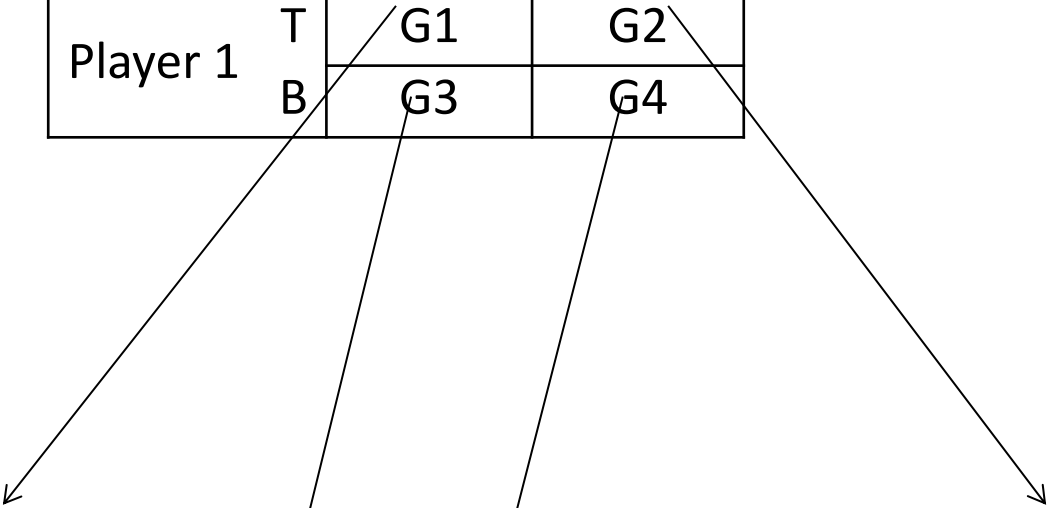
Example: two stage game of imperfect information

- **Stage 1:** Players 1 and 2 move simultaneously taking, respectively, actions $a_1 \in A_1$ and $a_2 \in A_2$
- **Stage 2:** Players 3 and 4 observe (a_1, a_2) , then move simultaneously taking, respectively, actions $a_3 \in A_3$ and $a_4 \in A_4$
- **Payoffs:** $u_i(a_1, a_2, a_3, a_4)$ for $i = 1, 2, 3, 4$
- **Solution:**
 - We solve the simultaneous - move game between players 3 and 4 in the second stage:
 - Players 1 and 2 anticipate the behaviour of players 3 and 4

example



		Player 2	
		L	R
Player 1	T	G1	G2
	B	G3	G4



G1		Player 4	
		L	R
Player 3	T	2, 2, <u>1</u> , 0	1, 1, <u>2</u> , <u>1</u>
	B	3, 3, 0, <u>0</u>	4, 4, 0, <u>0</u>

G2		Player 4	
		L	R
Player 3	T	3, 3, <u>1</u> , 0	0, 2, <u>1</u> , <u>1</u>
	B	4, 4, 0, 0	2, 2, 0, <u>1</u>

G3		Player 4	
		L	R
Player 3	T	3, 3, <u>1</u> , <u>2</u>	2, 2, <u>1</u> , 0
	B	0, 0, 0, <u>1</u>	0, 0, 0, 0

G4		Player 4	
		L	R
Player 3	T	2, 2, 0, <u>1</u>	2, 2, 0, 0
	B	2, 0, <u>1</u> , <u>1</u>	2, 2, <u>1</u> , 0

		Player 2	
		L	R
Player 1	T	1,1, <u>2</u> , 1	0, <u>2</u> , 1, 1
	B	<u>3</u> , <u>3</u> , 1, <u>2</u>	<u>2</u> , 0, 1, 1

G1

		Player 4	
		L	R
Player 3	T	2, 2, <u>1</u> , 0	<u>1</u> , <u>1</u> , <u>2</u> , <u>1</u>
	B	3, 3, 0, <u>0</u>	4, 4, 0, <u>0</u>

G2

		Player 4	
		L	R
Player 3	T	3, 3, <u>1</u> , 0	<u>0</u> , <u>2</u> , <u>1</u> , <u>1</u>
	B	4, 4, 0, 0	2, 2, 0, <u>1</u>

G3

		Player 4	
		L	R
Player 3	T	<u>3</u> , <u>3</u> , <u>1</u> , <u>2</u>	2, 2, <u>1</u> , 0
	B	0, 0, 0, <u>1</u>	0, 0, 0, 0

G4

		Player 4	
		L	R
Player 3	T	<u>2</u> , <u>2</u> , 0, <u>1</u>	2, 2, 0, 0
	B	<u>2</u> , 0, <u>1</u> , <u>1</u>	2, 2, <u>1</u> , 0

Backward induction outcome:

(B, L, T, L)

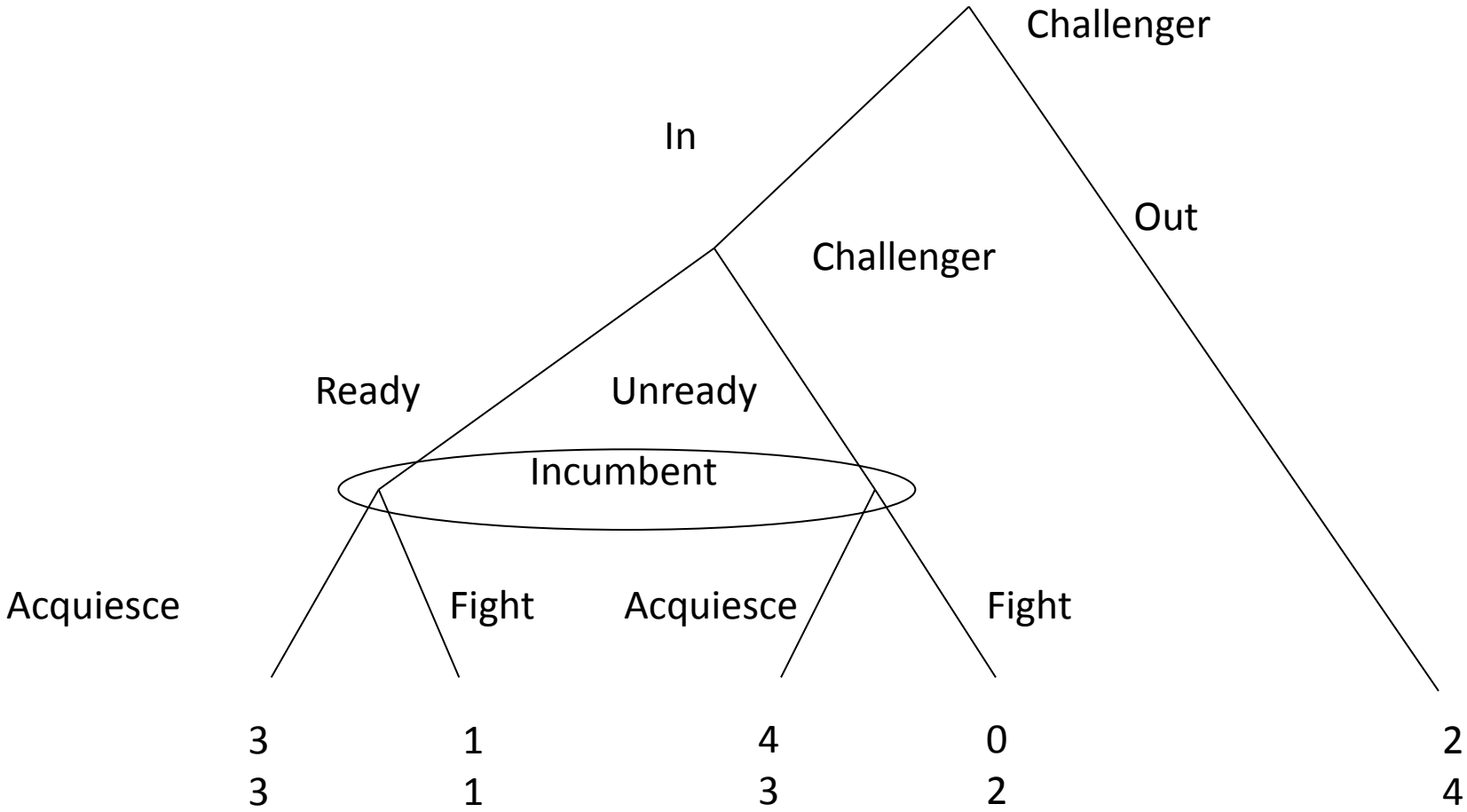
Subgame perfect Nash equilibrium

(B, L, (T, T, T, B), (R, R, L, L))

Example 2

Challenger's strategies: {(Out Ready), (Out Unready) (In ready), (In Unready)}

Incumbent's strategies: Acquiesce, Fight



		Incumbent	
		Acquiesce	Fight
Challenger	Out Ready	2, <u>4</u>	<u>2</u> , <u>4</u>
	Out Unready	2, <u>4</u>	<u>2</u> , <u>4</u>
	In Ready	3, <u>3</u>	1, 1
	In Unready	<u>4</u> , <u>3</u>	0, 2

Three Nash equilibria:
 (Out Ready, Fight);
 (Out Unready, Fight)
 (In unready, Acquiesce)

Consider the subgame starting in the decision node after Challenger's choice *In*

		Incumbent	
		Acquiesce	Fight
Challenger	Ready	3, <u>3</u>	<u>1</u> , 1
	Unready	<u>4</u> , <u>3</u>	0, 2

An unique Nash equilibrium: Unready, Acquiesce
 Then, only (In unready, Acquiesce) is **subgame perfect NE**

Example 3

		Incumbent	
		Acquiesce	Fight
Challenger	Ready	3, <u>3</u>	1, 1
	Unready	<u>4</u> , <u>3</u>	0, 2
	Out	2, <u>4</u>	<u>2</u> , <u>4</u>

Two Nash equilibria:
 (Out, Fight)
 (Unready, Acquiesce)

Both are SPNE

