

- Transmission of classical information through a quantum channel

• Alice wants to transmit messages $w \in \mathcal{W} = \{1, 2, \dots, 2^{nR}\}$ to Bob through the action of a memoryless quantum channel Φ .

• She encodes $w \in \mathcal{W}$ into an m -qubit state $\rho_w^{(n)}$ that is turned into $\sigma_w^{(n)} := \underbrace{\Phi \otimes \Phi \otimes \dots \otimes \Phi}_{n \text{ times}} [\rho_w^{(n)}]$ by n -uses of the channel.

• Bob tries to retrieve w by performing a POVM measurement $\{E_w\}_{w \in \mathcal{W}}$

$$\sum_{w \in \mathcal{W}} E_w^\dagger E_w = 1$$

• The probability of inferring w correctly is $\text{Tr}(\sigma_w^{(n)} E_w)$

• The average probability of error is then

$$P_{\text{avg}} = \frac{1}{2^{nR}} \sum_{w \in \mathcal{W}} (1 - \text{Tr}(\sigma_w^{(n)} E_w))$$

Definition 2.4.4.

158

The transmission of classical information at rate R through the quantum channel Φ is reliable if $p_{ew}^{(n)} \xrightarrow{n} 0$.

The rate R is achievable if there exist a $(2^{nR}, n)$ quantum encoding and POVM's $\{E_w^{(n)}\}_{w \in \mathcal{W}}$ such that R is reliable.

The capacity of Φ is the maximum reliable rate, the maximum being taken over all possible quantum encodings of $w \in \mathcal{W}$.

Remark: results are available only when the encoding is via product states: $w \rightarrow \rho_w^{(n)} = \rho_{w_1} \otimes \rho_{w_2} \otimes \dots \otimes \rho_{w_n}$;
the capacity of the quantum channel is then called product state capacity and denoted by $C^{(1)}(\Phi)$.

The product state capacity of a quantum memoryless channel Φ is given by $C^{(1)}(\Phi) = \chi(\Phi)$, where

$$\chi(\Phi) := \max_{\{p_x, \rho_x\}} \chi(\{p_x, \Phi[\rho_x]\}),$$

$$\chi(\{p_x, \Phi[\rho_x]\}) := S(\Phi[\sum_x p_x \rho_x]) - \sum_x p_x S(\Phi[\rho_x]),$$

the maximum being evaluated over all possible statistical ensembles

$$\{p_x, \rho_x\}_{x \in \mathcal{X}}, \quad p_x \geq 0; \quad \sum_{x \in \mathcal{X}} p_x = 1.$$

Remark

If $\rho(x) \in M_d(\mathbb{C}) \forall x \in \mathcal{X}$ then one can vary over ensembles consisting of at most d^2 pure states. (Try to prove it)

Proof that $R > \chi(\Phi)$ implies $p_{\text{aw}}^{(n)} \not\rightarrow 0$ as $n \rightarrow \infty$.

- Alice's encoding: $\{1, 2, \dots, 2^{nR}\} \ni w \mapsto \boxed{\rho_w^{(n)} = \rho_1(w) \otimes \dots \otimes \rho_n(w)}$
 $\rho_i(w) \in M_d(\mathbb{C})$, $\rho_w^{(n)} \in M_d(\mathbb{C})^{\otimes n} = \underbrace{M_d(\mathbb{C}) \otimes \dots \otimes M_d(\mathbb{C})}_{n \text{ times}}$

- Channel action: $\rho_w^{(n)} \mapsto \boxed{\sigma_w^{(n)} = \Phi^{\otimes n}[\rho_w^{(n)}] = \sigma_1(w) \otimes \dots \otimes \sigma_n(w)}$
 $\sigma_j(w) = \Phi[\rho_j(w)]$

- Bob's measurement: $\boxed{E := \{E_w\}_{w=0, 1, \dots, 2^{nR}}}$: $\boxed{\sum_{w=0}^{2^{nR}} E_w^\dagger E_w = 1}$

- Average error: $p_{\text{aw}}^{(n)} = \frac{\sum_{w=1}^{2^{nR}} (1 - \text{Tr}(\sigma_w^{(n)} E_w^\dagger E_w))}{2^{nR}}$

• Holevo's bound: $p, p_x \in M_d(\mathbb{C}) \Rightarrow \boxed{I_{\text{acc}}(X) \leq S(p) - \sum_{x \in \mathcal{X}} p(x) S(p_x) \leq \log d}$

Therefore: $\log(2^{nR}) \leq \log(d^n) \Rightarrow \boxed{R \leq \log d}$

• Fano's inequality: $\boxed{H(W|E) \leq h(p_{\text{aw}}^{(n)}) + p_{\text{aw}}^{(n)} \log(d^n)}$

$$h(p_{\text{aw}}^{(n)}) = -p_{\text{aw}}^{(n)} \log p_{\text{aw}}^{(n)} - (1-p_{\text{aw}}^{(n)}) \log(1-p_{\text{aw}}^{(n)})$$

$$H(W|E) = H(W) - I(W; E) = nR - I(W; E)$$

$$\geq nR - S\left(\sum_{w \in \mathcal{W}} \frac{1}{2^{nR}} \delta_w^{(n)}\right) + \sum_{w \in \mathcal{W}} \frac{1}{2^{nR}} S(\delta_w^{(n)}) \quad (\text{Holevo's bound})$$

$$\geq nR - S\left(\sum_{w \in \mathcal{W}} \frac{1}{2^{nR}} \delta_1(w) \otimes \dots \otimes \delta_n(w)\right) + \frac{1}{2^{nR}} \sum_{w \in \mathcal{W}} S(\delta_1(w) \otimes \dots \otimes \delta_n(w))$$

$$\geq nR - \sum_{j=1}^n \left(S\left(\sum_{w \in \mathcal{W}} \frac{1}{2^{nR}} \delta_j(w)\right) - \frac{1}{2^{nR}} \sum_{w \in \mathcal{W}} S(\delta_j(w)) \right) \quad (\text{sub-additivity})$$

$$\geq nR - n \chi(\mathbb{F}) \quad (\text{Holevo's bound})$$

Then,
$$P_{aw}^{(n)} n \log(d) \geq H(W|E) - h(P_{aw}^{(n)})$$

$$\geq n(R - \chi(\Phi))$$

$$\boxed{P_{aw}^{(n)} \geq \frac{R - \chi(\Phi)}{\log(d)} - \frac{h(P_{aw}^{(n)})}{n \log(d)}} \rightarrow 0 \quad \boxed{\Leftrightarrow R > \chi(\Phi)}$$

Proof that $R < \chi(\Phi)$ is an achievable rate.

• Idea: random coding,
$$\boxed{\text{Prob}(E) = \prod_{W \in \mathcal{X}} \prod_j P(x_j^E(w))}$$

$$W \ni w \xrightarrow{E} x_1^E(w) x_2^E(w) \dots x_n^E(w) =: x_E^{(n)} \in \underbrace{\mathcal{X} \times \mathcal{X} \times \dots \times \mathcal{X}}_{n \text{ times}}$$

$$\downarrow$$

$$P_E^{(n)}(w) = P(x_1^E(w)) \otimes \dots \otimes P(x_n^E(w)) \in M_d^{\otimes n}(\mathbb{C})$$

$$\downarrow$$

$$S_E^{(n)}(w) = \Phi^{\otimes n} [P_E^{(n)}(w)] = \underbrace{\Phi[x_1^E(w)]}_{\Phi[P(x_1^E(w))]} \otimes \dots \otimes \Phi[x_n^E(w)]$$

- $$\mathbb{E} [\rho_E^{(n)}(w)] = \sum_{\mathcal{E}} \text{Prob}(\mathcal{E}) \rho_E^{(n)}(w) = \bar{\rho}^{\otimes n} = \underbrace{\bar{\rho} \otimes \dots \otimes \bar{\rho}}_{n \text{ times}}$$

$$\bar{\rho} = \sum_{x \in \mathcal{X}} p(x) \rho(x) = \sum_{x \in \mathcal{X}} p(x) \mathbb{E}[p(x)]$$

- $\bar{\rho}^{\otimes n}$ typical subspace projector $\mathcal{P}_\epsilon^{(n)}$

$$\text{Tr}(\bar{\rho}^{\otimes n} (1 - \mathcal{P}_\epsilon^{(n)})) \leq \epsilon$$

- $$\rho(x) = \sum_j S_j(x) |S_j(x)\rangle \langle S_j(x)|$$

$$\left\{ \begin{aligned} \rho_E^{(n)}(w) &= \sum_{j^{(n)}} S_{j^{(n)}}^E(w) |S_{j^{(n)}}^E\rangle \langle S_{j^{(n)}}^E| \\ S_{j^{(n)}}^E(w) &= S_{j_1}(x_1^E(w)) \dots S_{j_n}(x_n^E(w)) \\ |S_{j^{(n)}}^E(w)\rangle &= |S_{j_1}(x_1^E(w))\rangle \otimes \dots \otimes |S_{j_n}(x_n^E(w))\rangle \end{aligned} \right.$$

$$\bar{S} = \sum_{x \in \mathcal{X}} p(x) S(\sigma(x))$$

$$T_{\mathcal{E}, \mathcal{E}}^{(n)}(w) := \text{LinSpan} \left\{ |S_{j^{(n)}}^{\mathcal{E}}(w)\rangle : \left| -\frac{1}{n} \log(S_{j^{(n)}}^{\mathcal{E}}(w)) - \bar{S} \right| \leq \mathcal{E} \right\}$$

$$P_{\mathcal{E}, \mathcal{E}}^{(n)}(w) : (\mathbb{C}^d)^{\otimes n} \rightarrow T_{\mathcal{E}, \mathcal{E}}^{(n)}(w) \quad \text{orthogonal projection}$$

$$- \sum_{\mathcal{E}} \text{Prob}(\mathcal{E}) \sum_{S_{j^{(n)}}^{\mathcal{E}}(w)} S_{j^{(n)}}^{\mathcal{E}}(w) \log S_{j^{(n)}}^{\mathcal{E}}(w) = \bar{S} \Rightarrow$$

$$\begin{aligned} \mathbb{E}[\text{Tr}(\sigma_w^{(n)} P_{\mathcal{E}, \mathcal{E}}^{(n)})] &\geq 1 - \mathcal{E} \\ \mathbb{E}[\text{Tr}(P_{\mathcal{E}, \mathcal{E}}^{(n)})] &\leq 2^{n(\bar{S} + \mathcal{E})} \end{aligned}$$

• Bob's square root decoding

$$E_w := P_{\mathcal{E}, \mathcal{E}}^{(n)}(w) P_{\mathcal{E}}^{(n)} \frac{1}{\sqrt{\sum_{\bar{w}} P_{\mathcal{E}}^{(n)} P_{\mathcal{E}, \mathcal{E}}^{(n)}(\bar{w}) P_{\mathcal{E}}^{(n)}}}$$

$$\sum_w P_{\mathcal{E}}^{(n)} P_{\mathcal{E}, \mathcal{E}}^{(n)}(w) P_{\mathcal{E}}^{(n)} =: \Pi_{\mathcal{E}, \mathcal{E}}^{(n)}$$

$$\sum_w E_w^\dagger E_w = \frac{1}{\sqrt{\Pi_{\mathcal{E}, \mathcal{E}}^{(n)}}} \left(\sum_w P_{\mathcal{E}}^{(n)} P_{\mathcal{E}, \mathcal{E}}^{(n)}(w) P_{\mathcal{E}}^{(n)} \right) \frac{1}{\sqrt{\Pi_{\mathcal{E}, \mathcal{E}}^{(n)}}} = \mathbb{1}$$

• Short-hand notation:

$$\left\{ \begin{array}{l} S_d(w) := S_{d^{(n)}}^E(w), \quad |S_{d^{(n)}}^E(w)\rangle := |S_d(w)\rangle \\ P_w := P_{E, E}^{(n)} = \left(\sum_d^1 |S_d(w)\rangle \langle S_d(w)| \right), \quad P := P_E^{(n)} \\ \Pi = \sum_w P P_w P, \quad |\tilde{S}_d(w)\rangle := P |S_d(w)\rangle \end{array} \right.$$

$$\begin{aligned} \text{Tr}(G_w^{(n)} E_w^\dagger E_w) &= \sum_R S_R(w) \langle \tilde{S}_R(w) | \frac{1}{\sqrt{\Pi}} \sum_e^1 | \tilde{S}_e(w) \rangle \langle \tilde{S}_e(w) | \frac{1}{\sqrt{\Pi}} | \tilde{S}_R(w) \rangle \\ &= \sum_R \sum_e^1 S_R(w) \left| \langle \tilde{S}_R(w) | \frac{1}{\sqrt{\Pi}} | \tilde{S}_e(w) \rangle \right|^2 \\ &\geq \sum_R^1 \sum_c^1 S_R(w) \left| \langle \tilde{S}_R(w) | \frac{1}{\sqrt{\Pi}} | \tilde{S}_c(w) \rangle \right|^2 \end{aligned}$$

$$\begin{aligned} P_{ev}^{(n)} &= \frac{1}{2^{nR}} \sum_w \left(1 - \text{Tr}(G_w^{(n)} E_w^\dagger E_w) \right) \leq \frac{1}{2^{nR}} \left(\sum_R^1 S_R(w) \left(1 - \sum_e^1 \left| \langle \tilde{S}_R(w) | \frac{1}{\sqrt{\Pi}} | \tilde{S}_e(w) \rangle \right|^2 \right) + \sum_R^{\text{II}} S_R(w) \right) \\ &= \frac{1}{2^{nR}} \sum_w \left(\sum_R^1 S_R(w) \left(1 - \sum_e^1 \left| \langle \tilde{S}_R(w) | \frac{1}{\sqrt{\Pi}} | \tilde{S}_e(w) \rangle \right|^2 \right) + \text{Tr}(G_w^{(n)} (1 - P_{E, E}^{(n)}(w))) \right) \end{aligned}$$

- $$P_{av}^{(n)} \leq \frac{1}{2^{NR}} \sum_w \left(\sum_R S_R(w) \left(1 - \left| \langle \tilde{S}_R(w) | \frac{1}{\sqrt{\pi}} | \tilde{S}_R(w) \rangle \right|^2 \right) + \text{Tr} \left(G_w^{(n)} (1 - P_{E,E}^{(n)}(w)) \right) \right)$$

- $$M_{Rw; R'w'} := \langle \tilde{S}_R(w) | \frac{1}{\sqrt{\pi}} | \tilde{S}_{R'}(w') \rangle ; \quad \pi = \sum_w \sum_R | \langle \tilde{S}_R(w) \rangle \langle \tilde{S}_R(w) |$$

$$\sum_{R', w'} M_{Rw; R'w'} M_{R'w'; R''w''} = \langle \tilde{S}_R(w) | \tilde{S}_{R''}(w'') \rangle$$

$$\langle \tilde{S}_R(w) | \tilde{S}_R(w) \rangle = \langle S_R(w) | P | S_R(w) \rangle \leq 1$$

||

$$\sum_{R', w'} M_{Rw; R'w'} M_{R'w'; R''w''} \geq (M_{Rw; R''w''})^2 = \left| \langle \tilde{S}_R(w) | \frac{1}{\sqrt{\pi}} | \tilde{S}_{R''}(w'') \rangle \right|^2$$

$$(1-x)^2 = (1-x)(1+x) \leq 2(1-x) \iff 0 \leq x \leq 1$$

- $$P_{av}^{(n)} \leq \frac{1}{2^{NR}} \sum_w \left(2 \sum_R S_R(w) \left(1 - \langle \tilde{S}_R(w) | \frac{1}{\sqrt{\pi}} | \tilde{S}_R(w) \rangle \right) + \text{Tr} \left(G_w^{(n)} (1 - P_{E,E}^{(n)}(w)) \right) \right)$$

• $M := [M_{Ew, E'w'}] : 2(1-M) = (1-M)^2 + 1-M^2 = (1-M^2)^2 \frac{1}{1+M^2} + 1-M^2$
 $\leq (1-M^2)^2 + 1-M^2 = 2-3M^2+M^4 \leq 2(1-M^2)+M^4$

• $P_{ew}^{(n)} \leq \frac{1}{2^{nR}} \sum_w \left(2 \sum_R^1 S_E(w) \left(1 - \langle \tilde{S}_R(w) | \tilde{S}_R(w) \rangle \right) + \sum_R^1 f_R(w) \sum_{R'}^1 \sum_{w'}^1 \left| \langle \tilde{S}_R(w) | \tilde{S}_{R'}(w') \rangle \right|^2 \right. \\ \left. + \text{Tr} \left(\zeta_w^{(n)} \left(1 - P_{\xi, E}^{(n)}(w) \right) \right) \right)$

$\leq \frac{1}{2^{nR}} \sum_w \left(2 \text{Tr} \left(\zeta_w^{(n)} \left(1 - P_{\xi, E}^{(n)} \right) \right) + \text{Tr} \left(\zeta_w^{(n)} \left(1 - P_{\xi, E}^{(n)}(w) \right) \right) \right. \\ \left. + \sum_{R, R'}^1 S_E(w) \left| \langle \tilde{S}_R(w) | \tilde{S}_{R'}(w) \rangle \right|^2 \right. \\ \left. + \sum_{R, R'}^1 S_E(w) \sum_{w' \neq w} \left| \langle \tilde{S}_R(w) | \tilde{S}_{R'}(w') \rangle \right|^2 \right)$

$$\begin{aligned} \bullet \sum_{E'} | \langle \tilde{s}_E(w) | \tilde{s}_{E'}(w) \rangle |^2 &= \sum_{E'} \langle s_E(w) | P | s_{E'}(w) \rangle \langle s_{E'}(w) | P | s_E(w) \rangle \\ &= \sum_{E'} \langle s_E(w) | (1-P) | s_{E'}(w) \rangle \langle s_{E'}(w) | (1-P) | s_E(w) \rangle \\ &= \langle s_E(w) | (1-P) P_w (1-P) | s_E(w) \rangle \end{aligned}$$

$$\bullet \sum_{E'} \sum_{w' \neq w} | \langle \tilde{s}_E(w) | \tilde{s}_{E'}(w') \rangle |^2 = \sum_{w' \neq w} \langle s_E(w) | P P_{w'} P | s_E(w) \rangle$$

$$\begin{aligned} \bullet P_{\text{av}}^{(n)} &\leq \frac{1}{2^{nR}} \sum_w \left(2 \text{Tr} \left(G_w^{(n)} (1 - P_E^{(n)}) \right) + \text{Tr} \left(G_w^{(n)} (1 - P_{E,\bar{E}}^{(n)}(w)) \right) \right) \\ &\quad + \text{Tr} \left(G_w^{(n)} (1 - P_E^{(n)}) P_{E,\bar{E}}^{(n)}(w) (1 - P_E^{(n)}) \right) \quad \left(\leq \text{Tr} \left(G_w^{(n)} (1 - P_E^{(n)}) \right) \right) \\ &\quad + \sum_{\bar{w} \neq w} \text{Tr} \left(G_w^{(n)} \left(P_E^{(n)} P_{E,\bar{E}}^{(n)}(\bar{w}) P_E^{(n)} \right) \right) \end{aligned}$$

$$P_{\omega}^{(n)} \leq \frac{1}{2^{nR}} \sum_w \left(3 \operatorname{Tr} \left(\bar{G}_w^{(n)} (1 + P_{\varepsilon}^{(n)}) \right) + \operatorname{Tr} \left(\bar{G}_w^{(n)} (1 - P_{\varepsilon, \varepsilon}^{(n)}(w)) \right) + \sum_{\bar{w} \neq w} \operatorname{Tr} \left(\bar{G}_w^{(n)} \left(P_{\varepsilon}^{(n)} P_{\varepsilon, \varepsilon}^{(n)}(\bar{w}) P_{\varepsilon}^{(n)} \right) \right) \right)$$

• Expectation with respect to random coding

$$\begin{aligned} \mathbb{E} [P_{\omega}^{(n)}] &\leq 3 \operatorname{Tr} \left(\bar{G}^{\otimes n} (1 + P_{\varepsilon}^{(n)}) \right) + \mathbb{E} \left(\operatorname{Tr} \left(\bar{G}_1^{(n)} (1 - P_{\varepsilon, \varepsilon}^{(n)}(1)) \right) \right) + (2^{nR} - 1) \operatorname{Tr} \left(\bar{G}^{\otimes n} P_{\varepsilon}^{(n)} \mathbb{E} [P_{\varepsilon, \varepsilon}^{(n)}(1)] P_{\varepsilon}^{(n)} \right) \\ &\leq 4 \varepsilon + (2^{nR} - 1) \operatorname{Tr} \left(\bar{G}^{\otimes n} P_{\varepsilon}^{(n)} \mathbb{E} [P_{\varepsilon, \varepsilon}^{(n)}(1)] P_{\varepsilon}^{(n)} \right) \end{aligned}$$

$$P_{\varepsilon}^{(n)} \bar{G}^{\otimes n} P_{\varepsilon}^{(n)} \leq 2^{-n} (S(\bar{G}) - \varepsilon) ; \quad \operatorname{Tr} \mathbb{E} [P_{\varepsilon, \varepsilon}^{(n)}(1)] = \mathbb{E} [\operatorname{Tr} (P_{\varepsilon, \varepsilon}^{(n)}(1))] \leq 2^{n(\bar{S} + \varepsilon)}$$

$$\begin{aligned} \mathbb{E} [P_{\omega}^{(n)}] &\leq 4 \varepsilon + (2^{nR} - 1) 2^{-n} (S(\bar{G}) - \varepsilon) 2^{n(\bar{S} + \varepsilon)} \\ &\leq 4 \varepsilon + 2^{-n} (S(\bar{G}) - \bar{S} - R) \end{aligned}$$

$$\begin{aligned} \bar{G} &= \sum_{x \in \mathcal{X}} p(x) \Phi [p(x)] \\ \bar{S} &= \sum_{x \in \mathcal{X}} p(x) S(\Phi [p(x)]) \end{aligned}$$