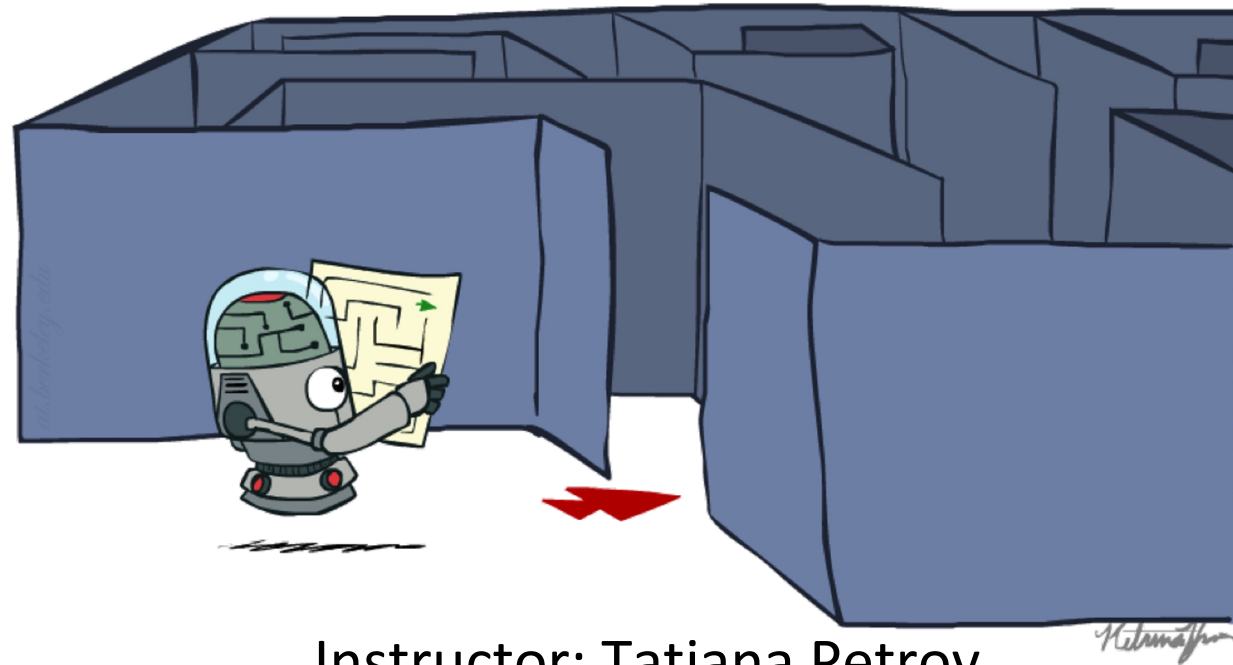


# Introduction to Artificial Intelligence

## Uninformed Search



Instructor: Tatjana Petrov

University of Trieste, Italy

# Uninformed Search

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No clue about how close a state is to the goal(s)

# Depth-First Search (DFS)

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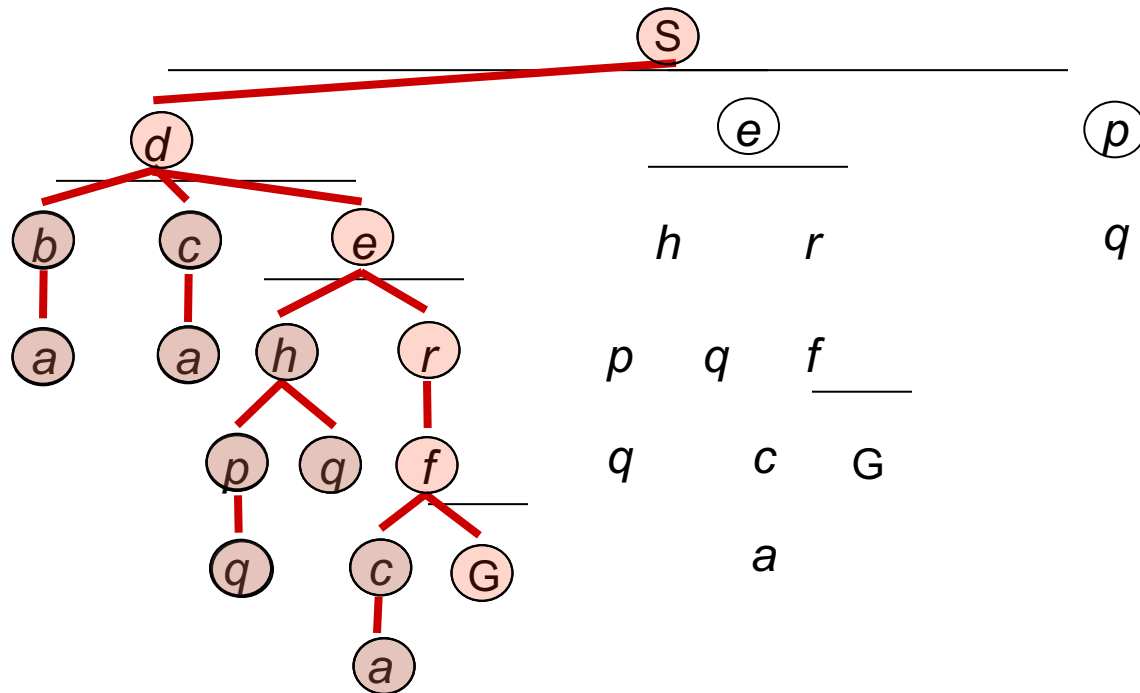
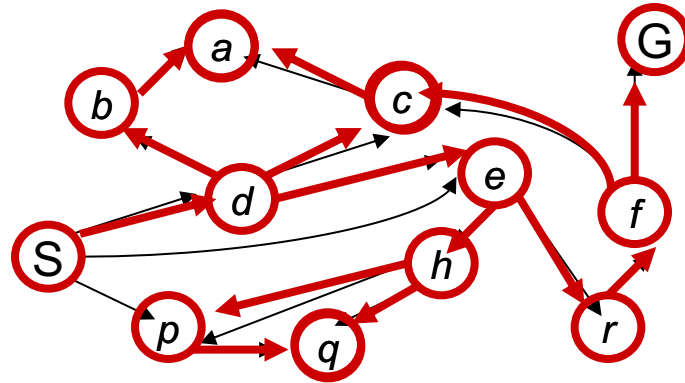


# Depth-First Search

Strategy: expand a deepest node first

Implementation:

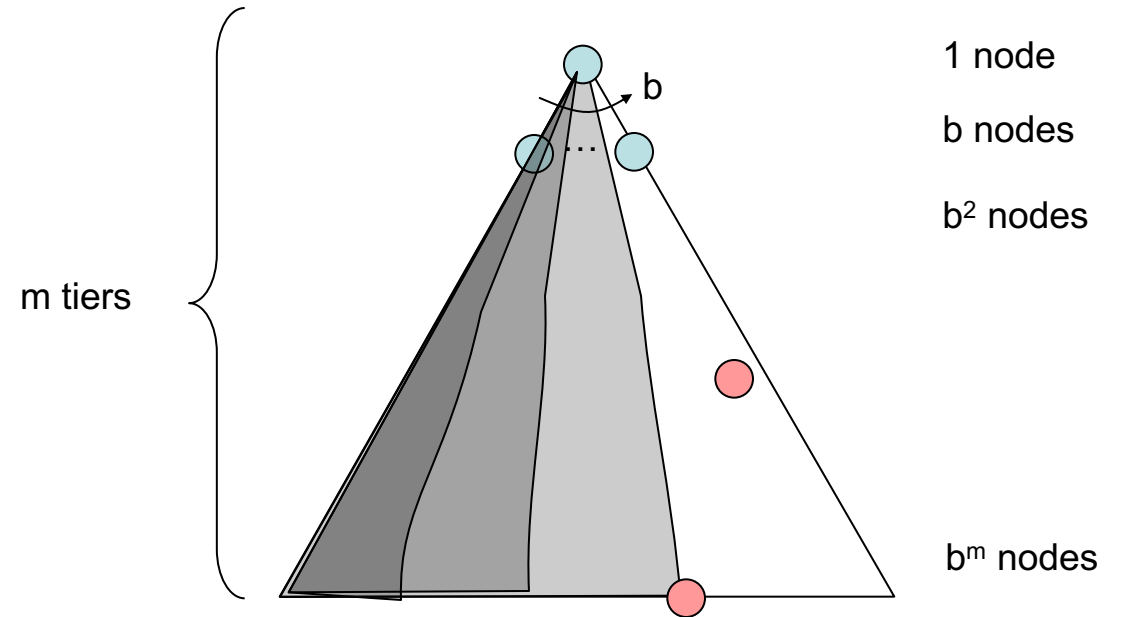
Fringe is a LIFO stack



~~sd~~  
~~se~~  
~~sp~~  
~~sdb~~  
~~sdc~~  
~~sde~~  
~~sdba~~  
~~sdc a~~  
~~sdc h~~  
~~sdc r~~

# Depth-First Search (DFS) Properties

- What nodes DFS expand?
  - Some left prefix of the tree.
  - Could process the whole tree!
  - If  $m$  is finite, takes time  $O(b^m)$
- How much space does the fringe take?
  - Only has siblings on path to root, so  $O(bm)$
- Is it complete?
  - $m$  could be infinite, so only if we prevent cycles (more later)
- Is it optimal?
  - No, it finds the “leftmost” solution, regardless of depth or cost



# Breadth-First Search (BFS)

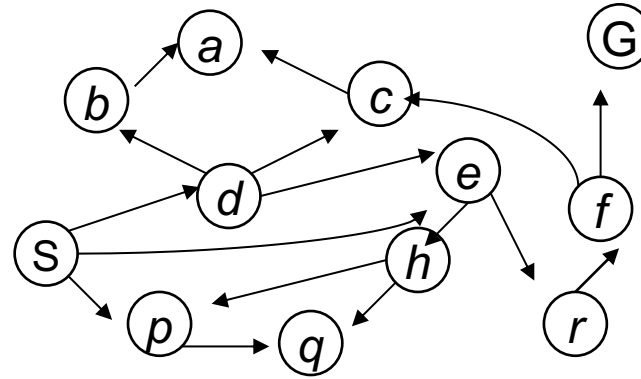
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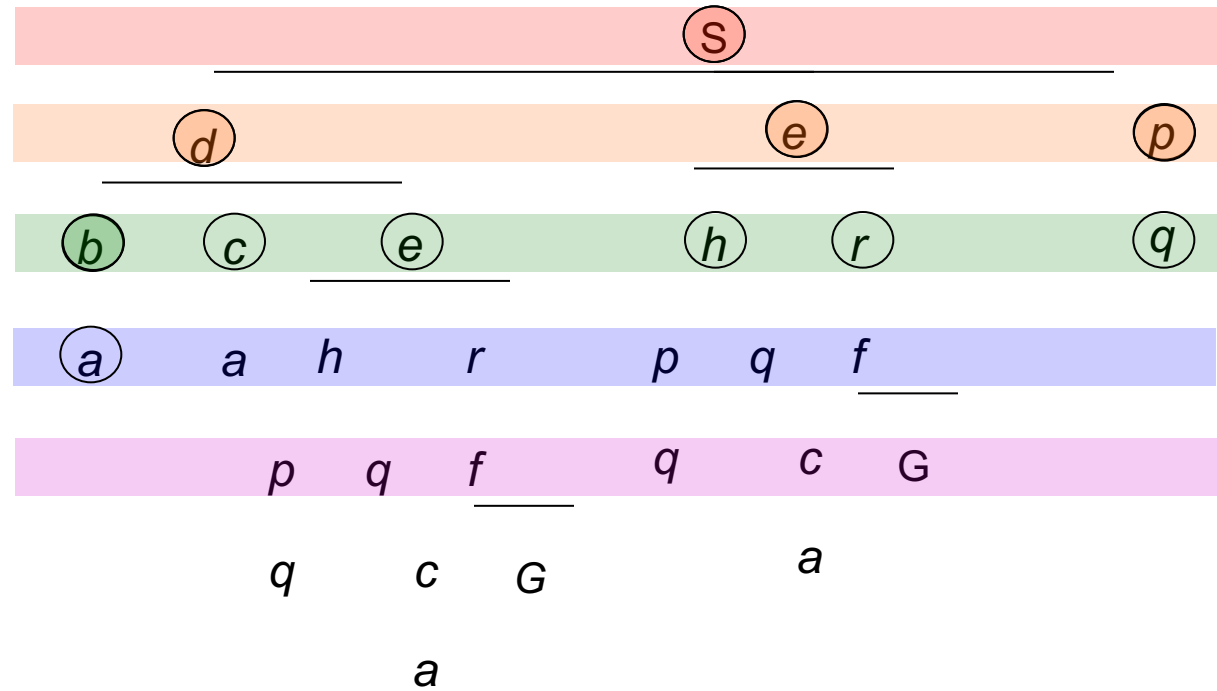
# Breadth-First Search

Strategy: expand a shallowest node first

Implementation: Fringe is a FIFO queue



Search Tiers

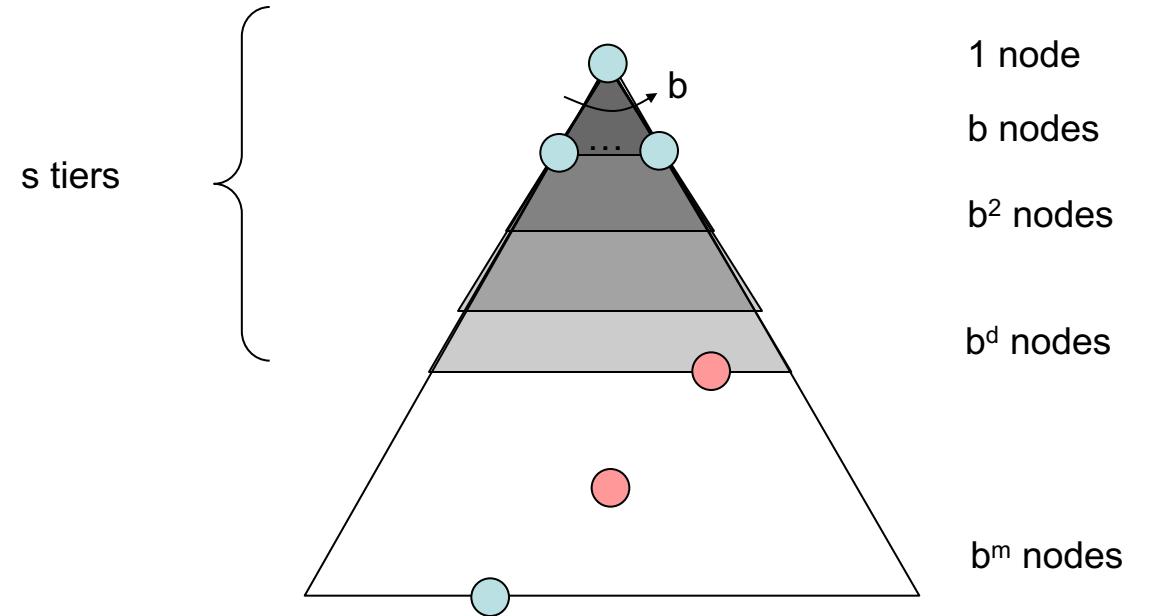


Handwritten notes in red:

- s
- sd
- se
- sp
- sdh
- sd c
- sd e
- seh
- ser
- spq

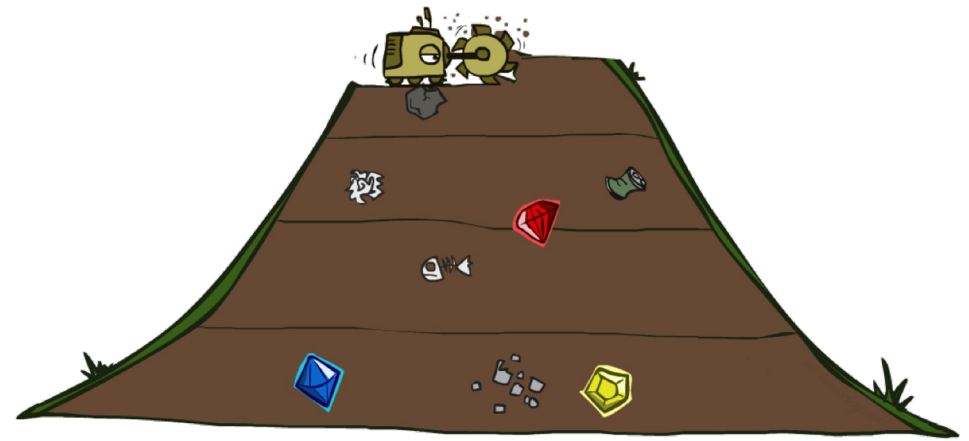
# Breadth-First Search (BFS) Properties

- What nodes does BFS expand?
  - Processes all nodes above shallowest solution
  - Let depth of shallowest solution be  $s$
  - Search takes time  $O(b^d)$
- How much space does the fringe take?
  - Has roughly the last tier, so  $O(b^d)$
- Is it complete?
  - $d$  must be finite if a solution exists, so yes!
- Is it optimal?
  - Only if costs are all 1 (more on costs later)





# Quiz: DFS vs BFS



# Quiz: DFS vs BFS

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- When will BFS outperform DFS?
- When will DFS outperform BFS?

# Video of Demo Maze Water DFS/BFS (part 1)

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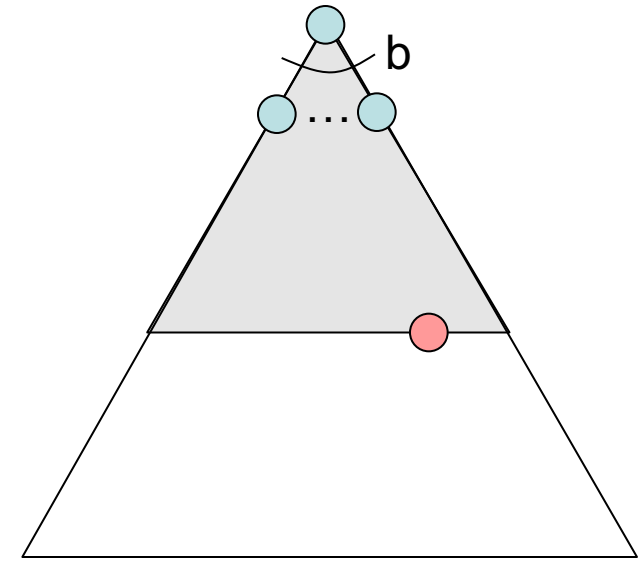
# Video of Demo Maze Water DFS/BFS (part 2)

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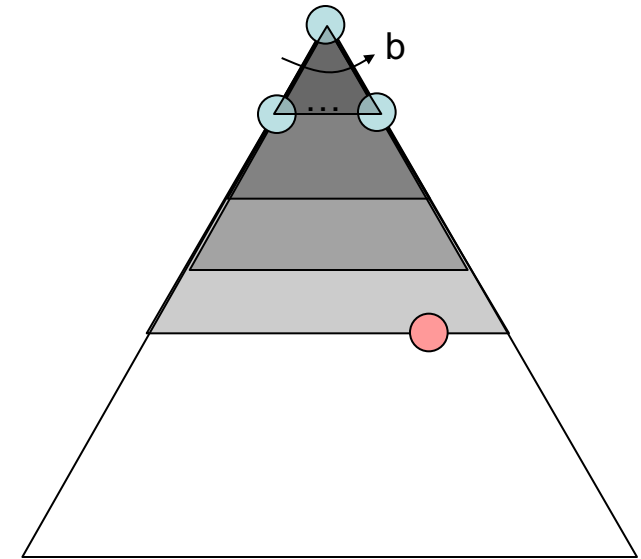
# Depth-limited Search (DLS)

- Idea: supply a depth limit  $\ell$  and treat all nodes at depth  $\ell$  as if they had no successors
- Time complexity?
  - Search takes time  $O(b^\ell)$
- Space complexity?
  - $O(b\ell)$
- Usefull when you know the **diameter** of the state-space graph

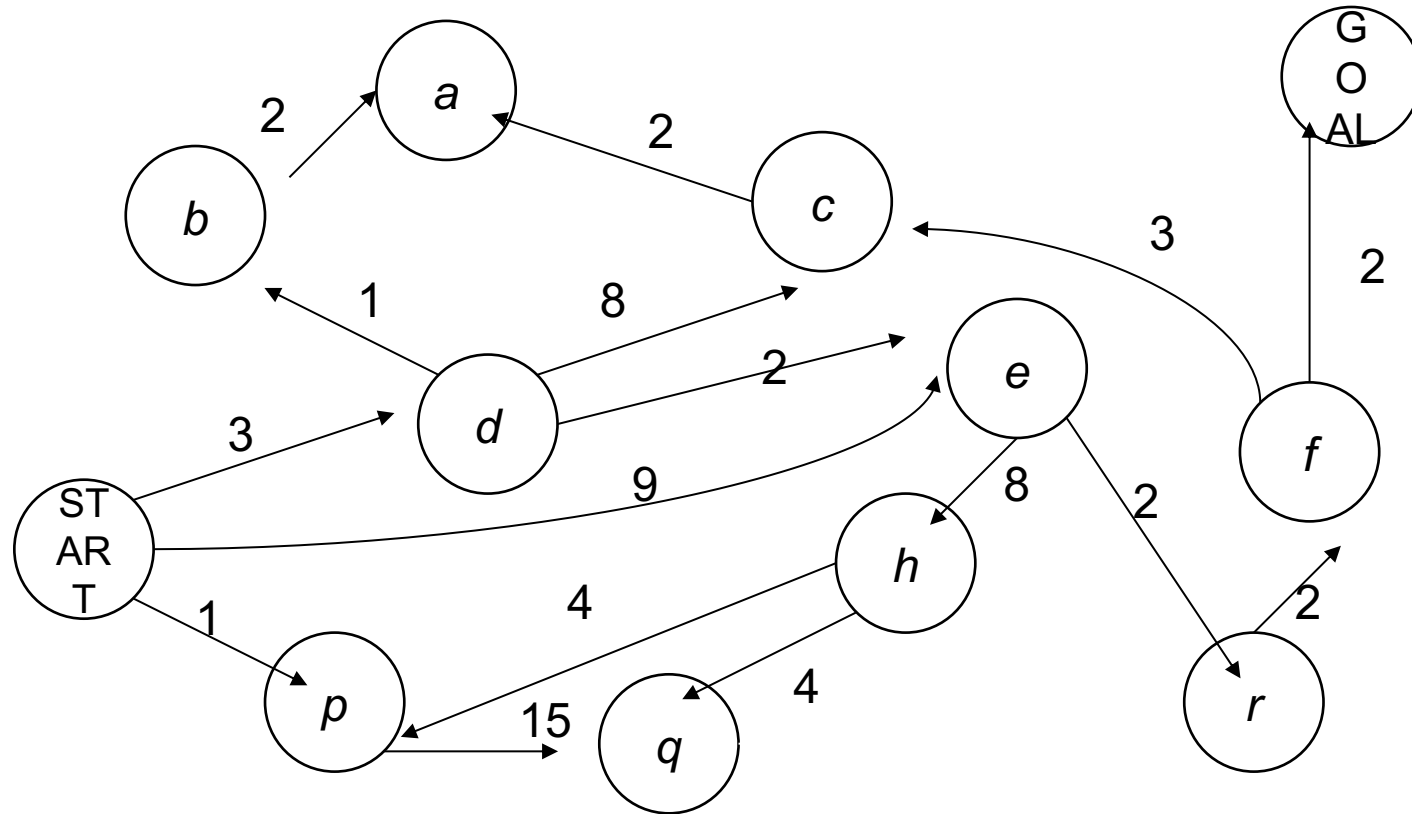


# Iterative Deepening Search (IDS)

- Idea: get DFS's space advantage with BFS's time / shallow-solution advantages
  - Run a DFS with depth limit 1. If no solution...
  - Run a DFS with depth limit 2. If no solution...
  - Run a DFS with depth limit 3. ....
- Time complexity?
  - Search takes time  $O(b^d)$
- Space complexity?
  - $O(bd)$
- Isn't that wastefully redundant?
  - Generally most work happens in the lowest level searched, so not so bad!



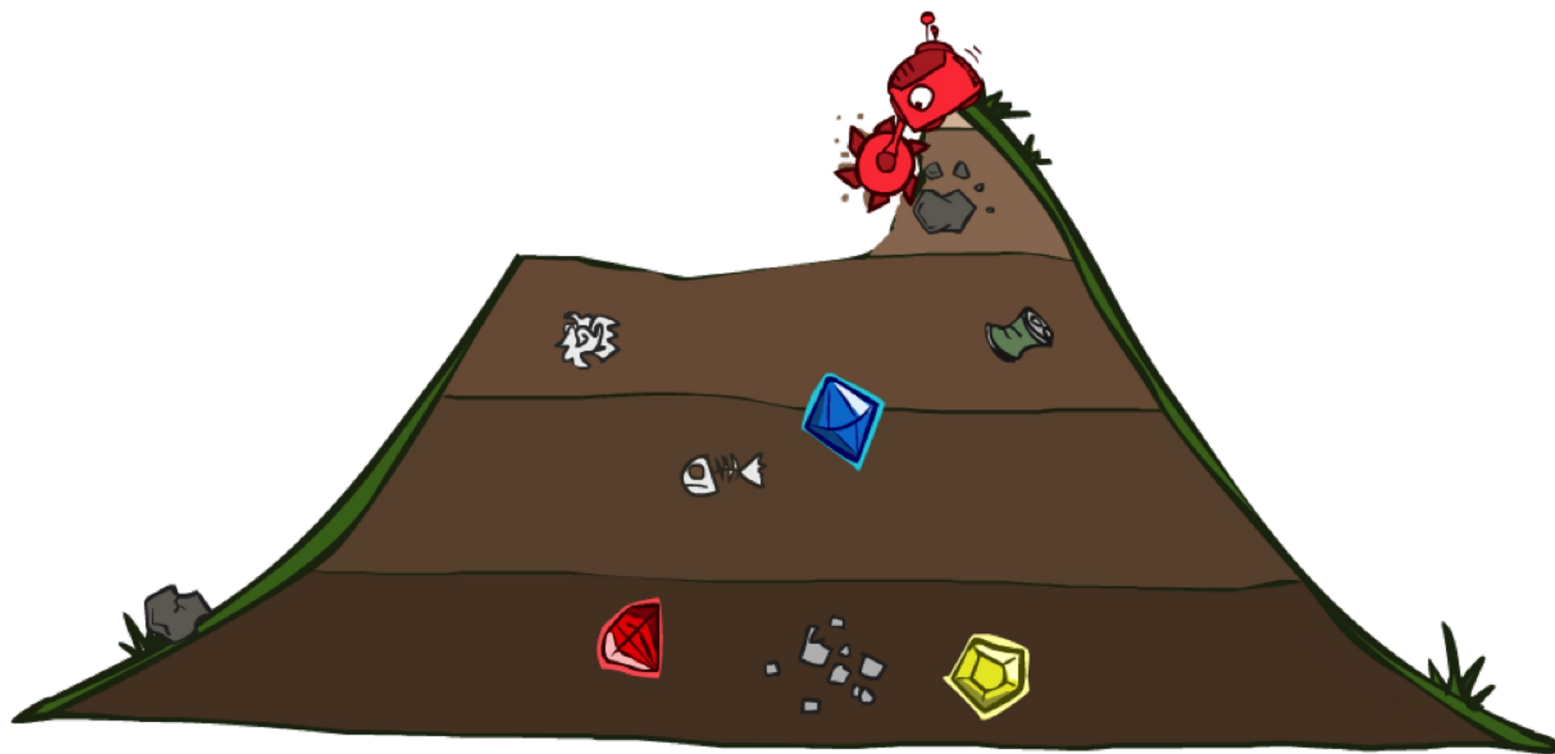
# Cost-Sensitive Search



BFS finds the shortest path in terms of number of actions. It does not find the least-cost path. We will now cover a similar algorithm which does find the least-cost path.

# Uniform Cost Search (UCS) (Dijkstra's algorithm)

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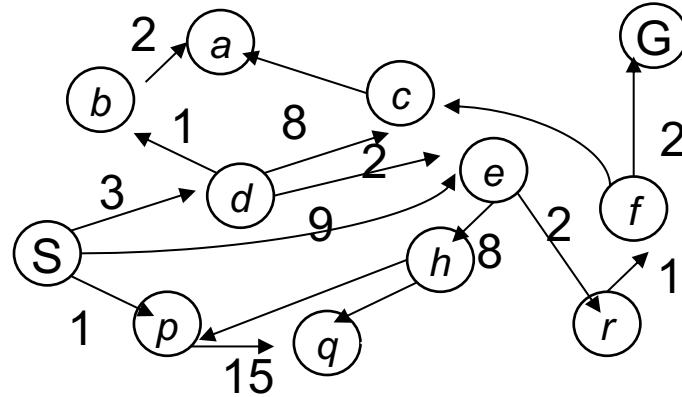




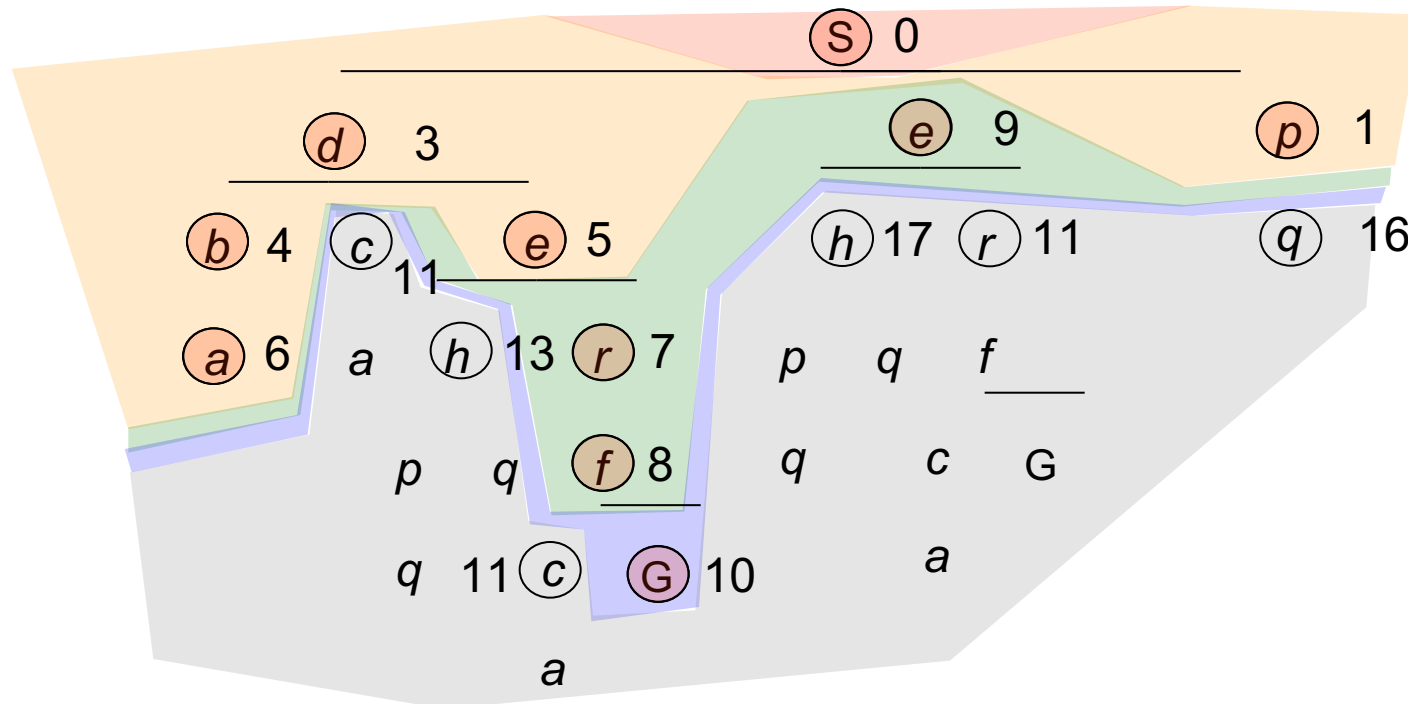
# Uniform Cost Search

Strategy: expand a cheapest node first:

Fringe is a priority queue (priority: cumulative cost)



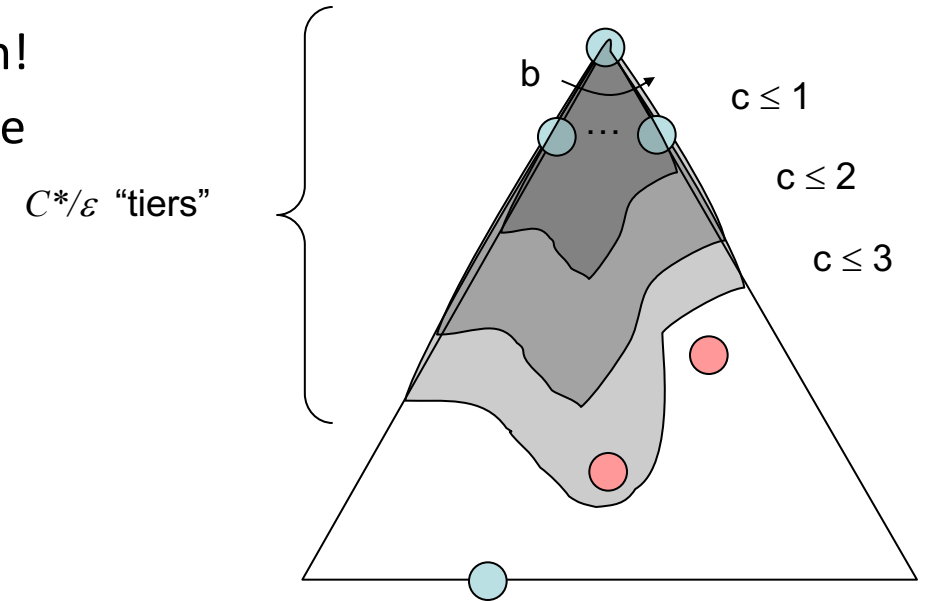
Cost contours



~~S~~  
~~sd 3~~  
~~se 9~~  
~~sp 1~~  
~~spq 16~~  
~~Sdb 4~~  
~~sdca 11~~  
~~sde 5~~  
~~Sdba 6~~  
~~sdch 13~~  
~~sder 7~~

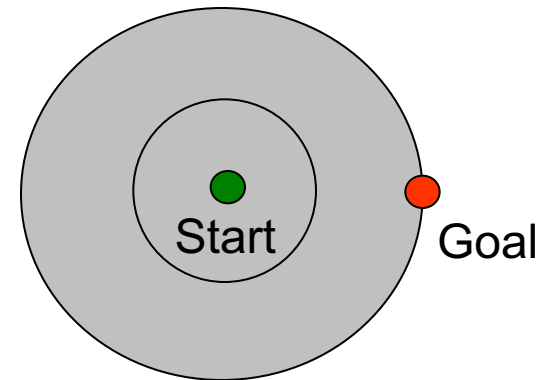
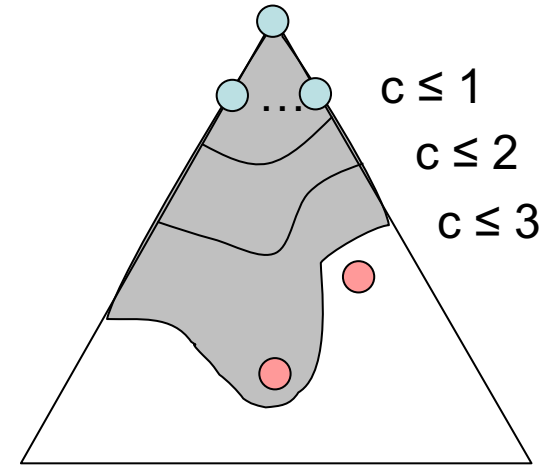
# Uniform Cost Search (UCS) Properties

- What nodes does UCS expand?
  - Processes all nodes with cost less than cheapest solution!
  - If that solution costs  $C^*$  and arcs cost at least  $\varepsilon$ , then the “effective depth” is roughly  $C^*/\varepsilon$
  - Takes time  $O(b^{I+C^*/\varepsilon})$  (exponential in effective depth)
- How much space does the fringe take?
  - Has roughly the last tier, so  $O(b^{I+C^*/\varepsilon})$
- Is it complete?
  - Assuming best solution has a finite cost and minimum arc cost is positive, yes!
- Is it optimal?
  - Yes! (Proof next lecture via  $A^*$ )



# Uniform Cost Issues

- Remember: UCS explores increasing cost contours
- The good: UCS is complete and optimal!
- The bad:
  - Explores options in every “direction”
  - No information about goal location
- We’ll fix that soon!



# Video of Demo Contours UCS Pacman Small Maze

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# Video of Demo Empty UCS

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# Video of Demo Maze with Deep/Shallow Water --- DFS, BFS, or UCS? (part 1)

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# Video of Demo Maze with Deep/Shallow Water --- DFS, BFS, or UCS? (part 2)

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# Video of Demo Maze with Deep/Shallow Water --- DFS, BFS, or UCS? (part 3)

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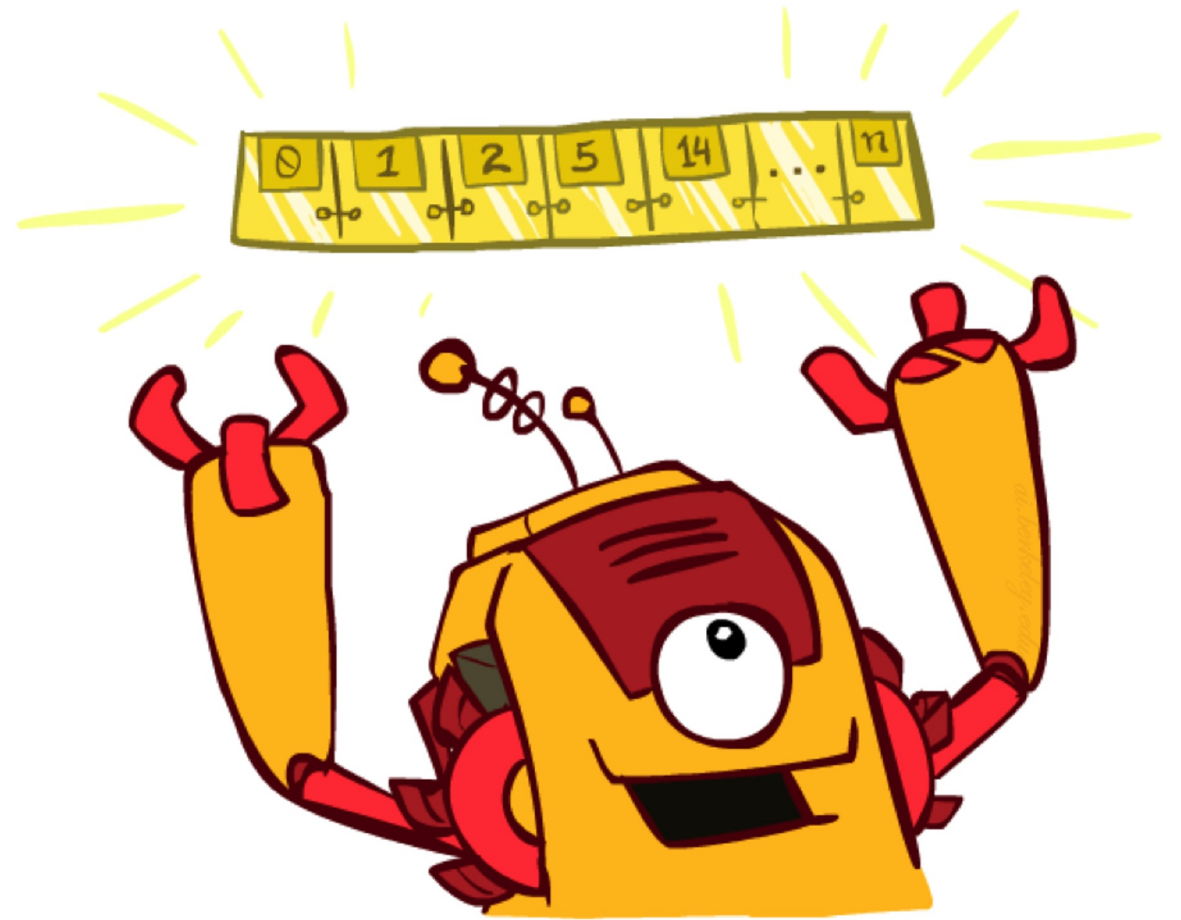


# Comparing uninformed search algorithms

| Criterion     | Breadth-First    | Uniform-Cost                          | Depth-First | Depth-Limited | Iterative Deepening |
|---------------|------------------|---------------------------------------|-------------|---------------|---------------------|
| Complete?     | Yes <sup>1</sup> | Yes <sup>1,2</sup>                    | No          | No            | Yes <sup>1</sup>    |
| Optimal cost? | Yes <sup>3</sup> | Yes                                   | No          | No            | Yes <sup>3</sup>    |
| Time          | $O(b^d)$         | $O(b^{1+\lceil C^*/\epsilon \rceil})$ | $O(b^m)$    | $O(b^\ell)$   | $O(b^d)$            |
| Space         | $O(b^d)$         | $O(b^{1+\lceil C^*/\epsilon \rceil})$ | $O(bm)$     | $O(b\ell)$    | $O(bd)$             |

# The One Queue

- All these search algorithms are the same except for fringe strategies
  - Conceptually, all fringes are priority queues (i.e. collections of nodes with attached priorities)
  - Practically, for DFS and BFS, you can avoid the  $\log(n)$  overhead from an actual priority queue, by using stacks and queues
  - Can even code one implementation that takes a variable queuing object



# Search and Models

- Search operates over models of the world
  - The agent doesn't actually try all the plans out in the real world!
  - Planning is all “in simulation”
  - Your search is only as good as your models...

