

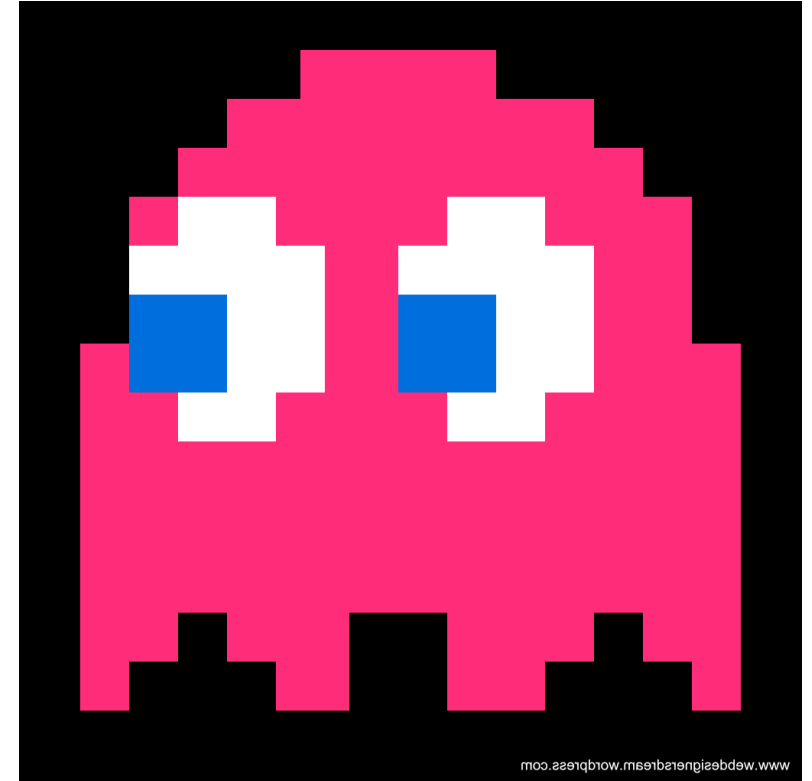
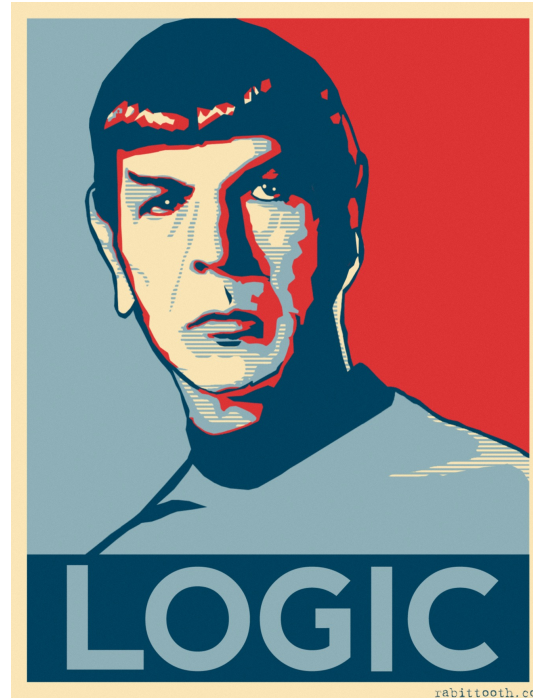
# Announcements

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- HW
- Project

# 272SM: Artificial Intelligence

## Propositional Logic I



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# Outline

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## 1. Propositional Logic I

- Basic concepts of knowledge, logic, reasoning
- Propositional logic: syntax and semantics, Pacworld example

## 2. Propositional logic II

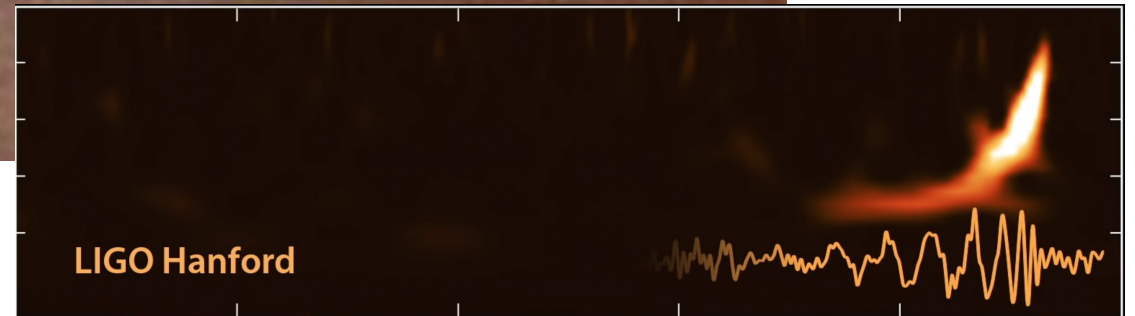
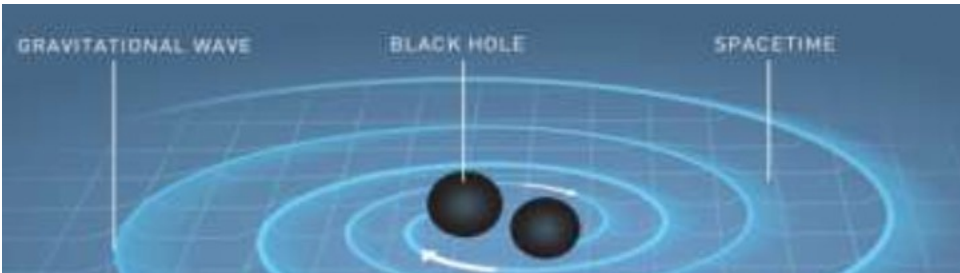
- Inference by theorem proving (briefly) and model checking
- A Pac agent using propositional logic

# Agents that know things

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- Agents acquire knowledge through perception, learning, language
  - Knowledge of the effects of actions (“transition model”)
  - Knowledge of how the world affects sensors (“sensor model”)
  - Knowledge of the current state of the world
- Can keep track of a partially observable world
- Can formulate plans to achieve goals
- Can design and build gravitational wave detectors.....

# LIGO



# Knowledge, contd.

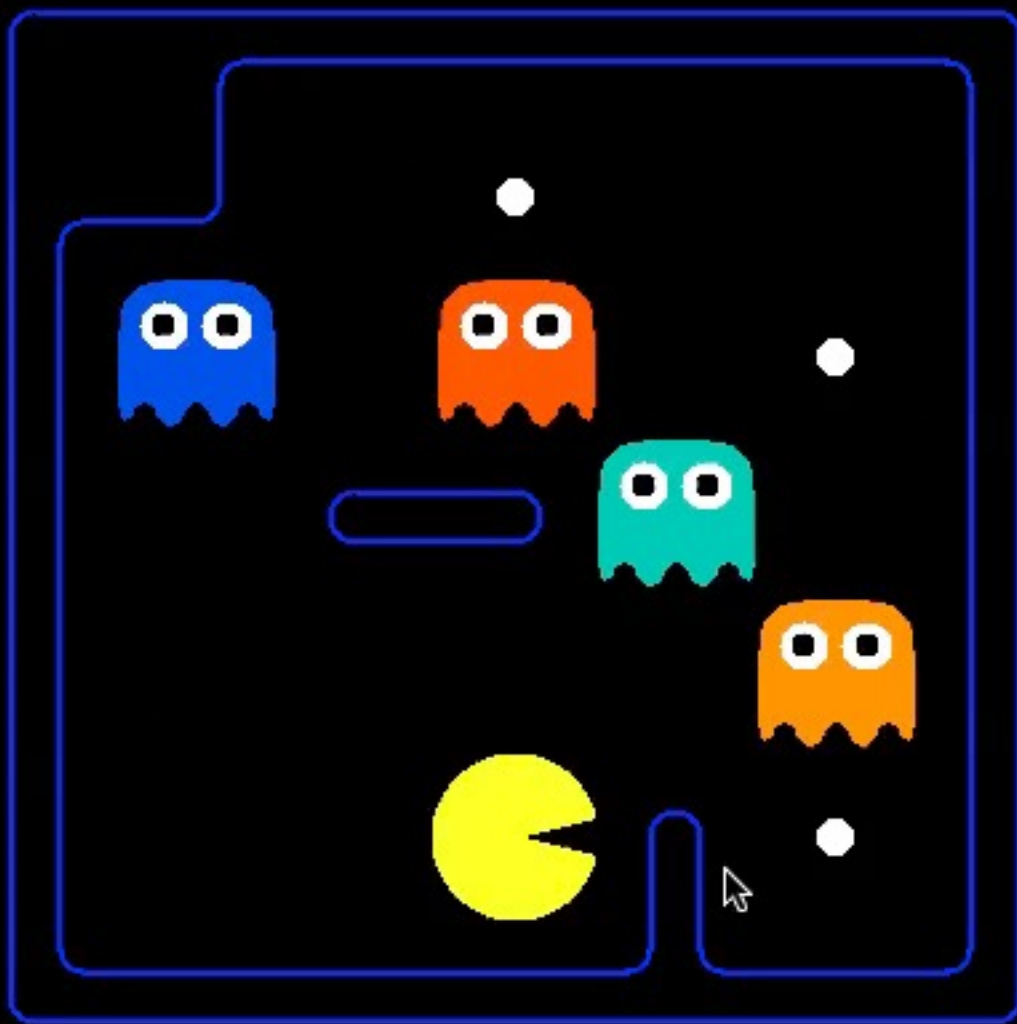
- Knowledge base = set of sentences in a formal language
- Declarative approach to building an agent (or other system):
  - **Tell** it what it needs to know (or have it **Learn** the knowledge)
  - Then it can **Ask** itself what to do—answers should follow from the KB
- Agents can be viewed at the **knowledge level**  
i.e., what they **know**, regardless of how implemented
- A single inference algorithm can answer any answerable question

Knowledge base

Domain-specific facts

Inference engine

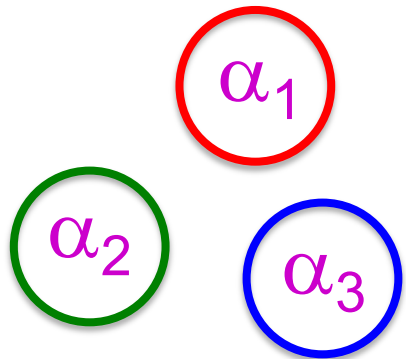
Generic code



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# Logic

- **Syntax:** What sentences are allowed?
- **Semantics:**
  - What are the **possible worlds**?
  - Which sentences are **true** in which worlds? (i.e., **definition** of truth)



*Syntaxland*



*Semanticsland*



# Different kinds of logic

## ■ Propositional logic

- Syntax:  $P \vee (\neg Q \wedge R)$ ;  $X_1 \Leftrightarrow (\text{Raining} \Rightarrow \neg \text{Sunny})$
- Possible world:  $\{P=\text{true}, Q=\text{true}, R=\text{false}, S=\text{true}\}$  or 1101
- Semantics:  $\alpha \wedge \beta$  is true in a world iff  $\alpha$  is true and  $\beta$  is true (etc.)

## ■ First-order logic

- Syntax:  $\forall x \exists y P(x,y) \wedge \neg Q(\text{Joe}, f(x)) \Rightarrow f(x)=f(y)$
- Possible world: Objects  $o_1, o_2, o_3$ ;  $P$  holds for  $\langle o_1, o_2 \rangle$ ;  $Q$  holds for  $\langle o_3 \rangle$ ;  $f(o_1)=o_1$ ;  $\text{Joe}=o_3$ ; etc.
- Semantics:  $\phi(\sigma)$  is true in a world if  $\sigma=o_j$  and  $\phi$  holds for  $o_j$ ; etc.

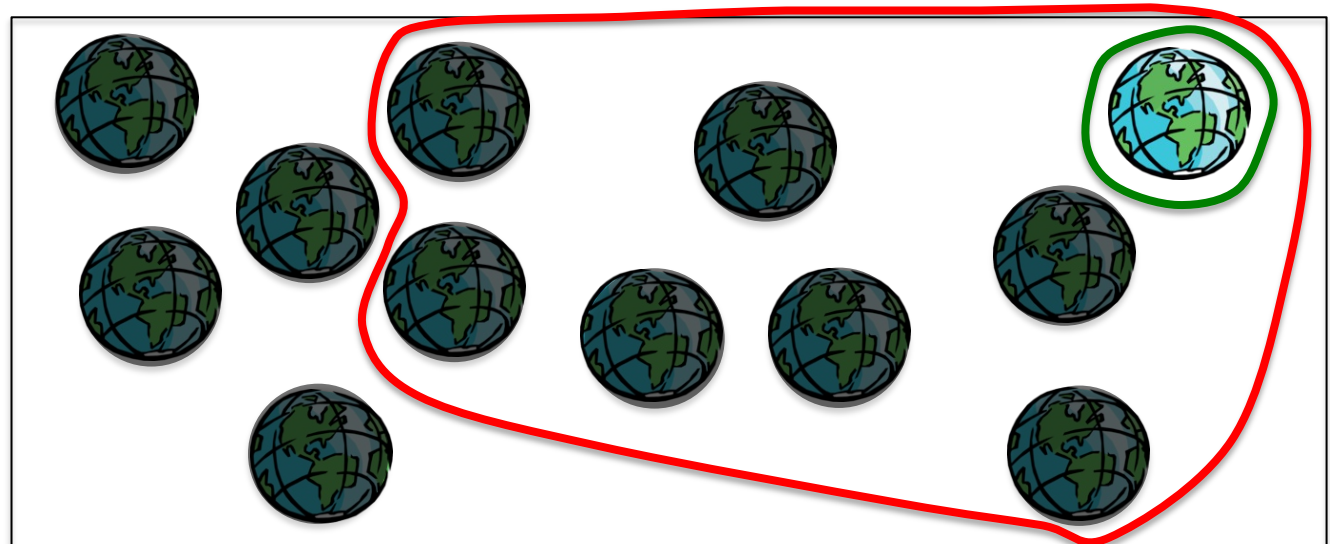
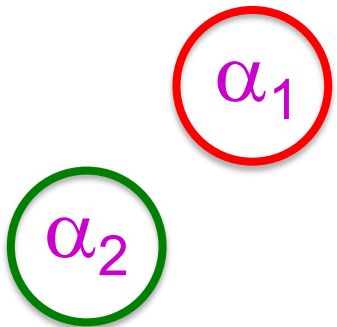
# Different kinds of logic, contd.

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- Relational databases:
  - Syntax: ground relational sentences, e.g., *Sibling(Ali,Bo)*
  - Possible worlds: (typed) objects and (typed) relations
  - Semantics: sentences in the DB are true, everything else is false
    - Cannot express disjunction, implication, universals, etc.
    - Query language (SQL etc.) typically some variant of first-order logic
    - Often augmented by first-order rule languages, e.g., Datalog
  - Knowledge graphs (roughly: relational DB + ontology of types and relations)
    - Google Knowledge Graph: 5 billion entities, 500 billion facts, >30% of queries
    - Facebook network: 2.93 billion people, trillions of posts, maybe quadrillions of facts

# Inference: entailment

- **Entailment:**  $\alpha \models \beta$  (“ $\alpha$  entails  $\beta$ ” or “ $\beta$  follows from  $\alpha$ ”) iff in every world where  $\alpha$  is true,  $\beta$  is also true
  - I.e., the  $\alpha$ -worlds are a **subset** of the  $\beta$ -worlds [ $models(\alpha) \subseteq models(\beta)$ ]
- In the example,  $\alpha_2 \models \alpha_1$
- (Say  $\alpha_2$  is  $\neg Q \wedge R \wedge S \wedge W$   
 $\alpha_1$  is  $\neg Q$ )



# Inference: proofs

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- A proof is a *demonstration* of entailment between  $\alpha$  and  $\beta$
- *Sound* algorithm: everything it claims to prove is in fact entailed
- *Complete* algorithm: every that is entailed can be proved

# Inference: proofs

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- Method 1: *model-checking*
  - For every possible world, if  $\alpha$  is true make sure that is  $\beta$  true too
  - OK for propositional logic (finitely many worlds); not easy for first-order logic
- Method 2: *theorem-proving*
  - Search for a sequence of proof steps (applications of *inference rules*) leading from  $\alpha$  to  $\beta$
  - E.g., from  $P$  and  $(P \Rightarrow Q)$ , infer  $Q$  by *Modus Ponens*

# Propositional logic syntax

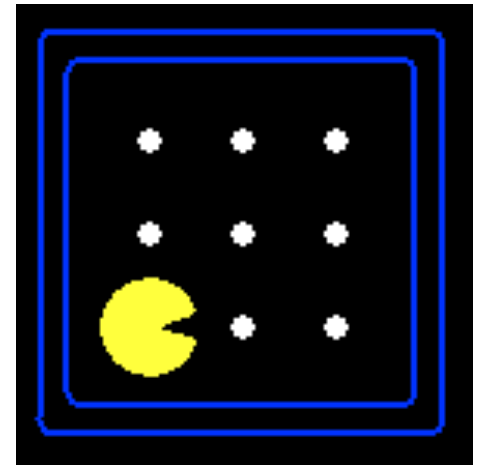
- Given: a set of proposition symbols  $\{X_1, X_2, \dots, X_n\}$ 
  - (we often add **True** and **False** for convenience)
- $X_i$  is a sentence
- If  $\alpha$  is a sentence then  $\neg\alpha$  is a sentence
- If  $\alpha$  and  $\beta$  are sentences then  $\alpha \wedge \beta$  is a sentence
- If  $\alpha$  and  $\beta$  are sentences then  $\alpha \vee \beta$  is a sentence
- If  $\alpha$  and  $\beta$  are sentences then  $\alpha \Rightarrow \beta$  is a sentence
- If  $\alpha$  and  $\beta$  are sentences then  $\alpha \Leftrightarrow \beta$  is a sentence
- And p.s. there are no other sentences!

# Propositional logic semantics

- Let  $m$  be a model assigning true or false to  $\{X_1, X_2, \dots, X_n\}$
- If  $\alpha$  is a symbol then its truth value is given in  $m$
- $\neg\alpha$  is true in  $m$  iff  $\alpha$  is false in  $m$
- $\alpha \wedge \beta$  is true in  $m$  iff  $\alpha$  is true in  $m$  and  $\beta$  is true in  $m$
- $\alpha \vee \beta$  is true in  $m$  iff  $\alpha$  is true in  $m$  or  $\beta$  is true in  $m$
- $\alpha \Rightarrow \beta$  is true in  $m$  iff  $\alpha$  is false in  $m$  or  $\beta$  is true in  $m$
- $\alpha \Leftrightarrow \beta$  is true in  $m$  iff  $\alpha \Rightarrow \beta$  is true in  $m$  and  $\beta \Rightarrow \alpha$  is true in  $m$

# Example: Partially observable Pacman

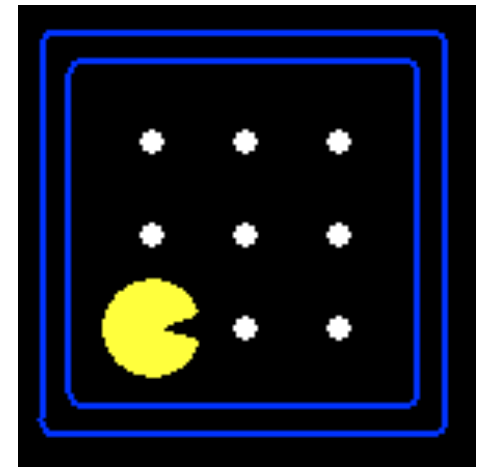
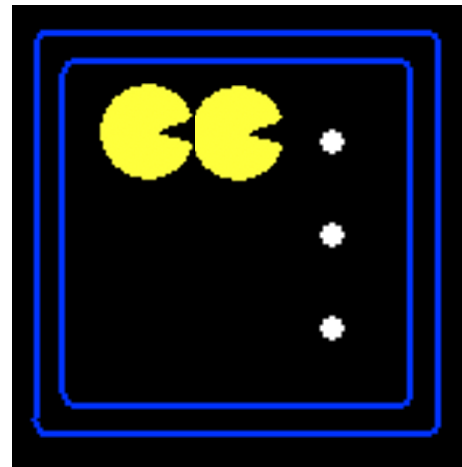
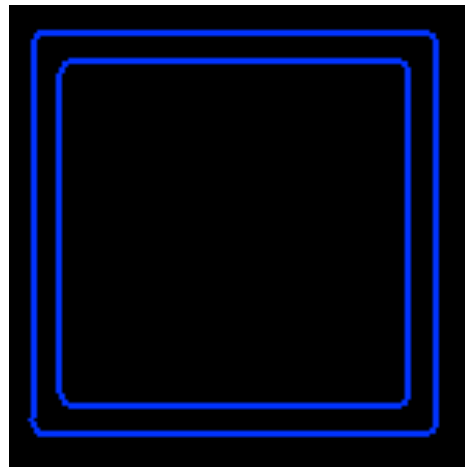
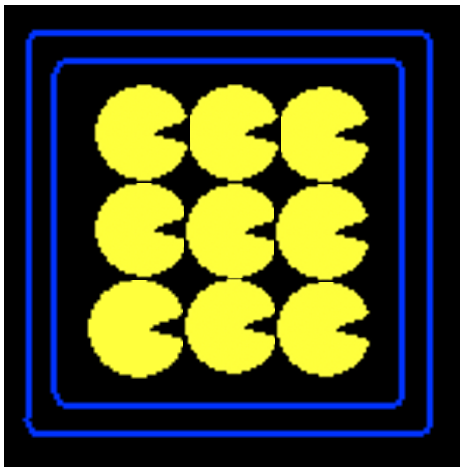
- Pacman knows the map but perceives just wall/gap to NSEW
- Formulation: *what variables do we need?*
  - Wall locations
    - $Wall_{0,0}$  there is a wall at [0,0]
    - $Wall_{0,1}$  there is a wall at [0,1], etc. ( $N$  symbols for  $N$  locations)
  - Percepts
    - ~~■  $Blocked_W$  (blocked by wall to my West) etc.~~
    - $Blocked_W_0$  (blocked by wall to my West at time 0) etc. ( $4T$  symbols for  $T$  time steps)
  - Actions
    - $W_0$  (Pacman moves West at time 0),  $E_0$  etc. ( $4T$  symbols)
  - Pacman's location
    - $At_{0,0}_0$  (Pacman is at [0,0] at time 0),  $At_{0,1}_0$  etc. ( $NT$  symbols)





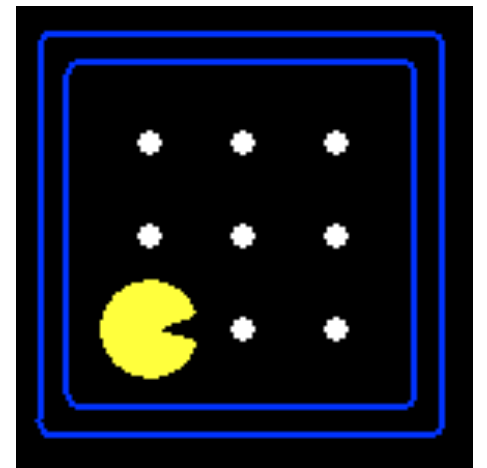
# How many possible worlds?

- $N$  locations,  $T$  time steps  $\Rightarrow N + 4T + 4T + NT = O(NT)$  variables
- $O(2^{NT})$  possible worlds!
- $N=200, T=400 \Rightarrow \sim 10^{24000}$  worlds
- Each world is a complete “history”
  - But most of them are pretty weird!



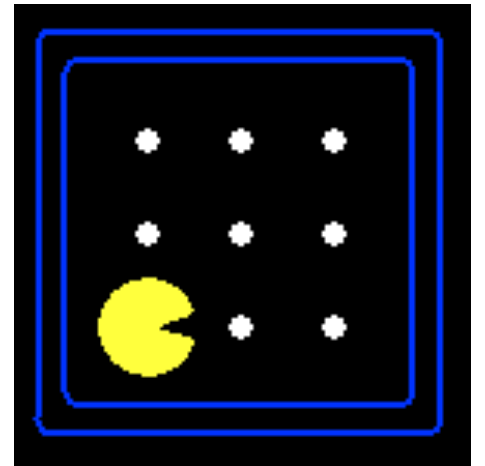
# Pacman's knowledge base: Map

- Pacman knows where the walls are:
  - $Wall_{0,0} \wedge Wall_{0,1} \wedge Wall_{0,2} \wedge Wall_{0,3} \wedge Wall_{0,4} \wedge Wall_{1,4} \wedge \dots$
- Pacman knows where the walls aren't!
  - $\neg Wall_{1,1} \wedge \neg Wall_{1,2} \wedge \neg Wall_{1,3} \wedge \neg Wall_{2,1} \wedge \neg Wall_{2,2} \wedge \dots$



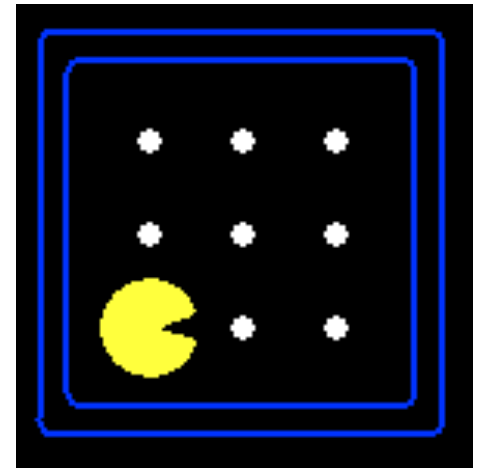
# Pacman's knowledge base: Initial state

- Pacman doesn't know where he is
- But he knows he's somewhere!
  - $At_{1,1}_0 \vee At_{1,2}_0 \vee At_{1,3}_0 \vee At_{2,1}_0 \vee \dots$



# Pacman's knowledge base: Sensor model

- State facts about how Pacman's percepts arise...
  - $\langle \text{Percept variable at } t \rangle \Leftrightarrow \langle \text{some condition on world at } t \rangle$
- Pacman perceives a wall to the West at time  $t$  **if and only if** he is in  $x,y$  and there is a wall at  $x-1,y$ 
  - $\text{Blocked\_W\_0} \Leftrightarrow ((\text{At\_1,1\_0} \wedge \text{Wall\_0,1}) \vee (\text{At\_1,2\_0} \wedge \text{Wall\_0,2}) \vee (\text{At\_1,3\_0} \wedge \text{Wall\_0,3}) \vee \dots)$
  - $4T$  sentences, each of size  $O(N)$
  - Note: these are valid for any map



# Pacman's knowledge base: Transition model

- How does each *state variable* at each time gets its value?
  - Here we care about location variables, e.g.,  $At_{3,3}_{17}$
- A state variable  $X$  gets its value according to a *successor-state axiom*
  - $X_t \Leftrightarrow [X_{t-1} \wedge \neg(\text{some action}_{t-1} \text{ made it false})] \vee [\neg X_{t-1} \wedge (\text{some action}_{t-1} \text{ made it true})]$
- For Pacman location:
  - $At_{3,3}_{17} \Leftrightarrow [At_{3,3}_{16} \wedge \neg((\neg Wall_{3,4} \wedge N_{16}) \vee (\neg Wall_{4,3} \wedge E_{16}) \vee \dots)] \vee [\neg At_{3,3}_{16} \wedge ((At_{3,2}_{16} \wedge \neg Wall_{3,3} \wedge N_{16}) \vee (At_{2,3}_{16} \wedge \neg Wall_{3,3} \wedge N_{16}) \vee \dots)]$

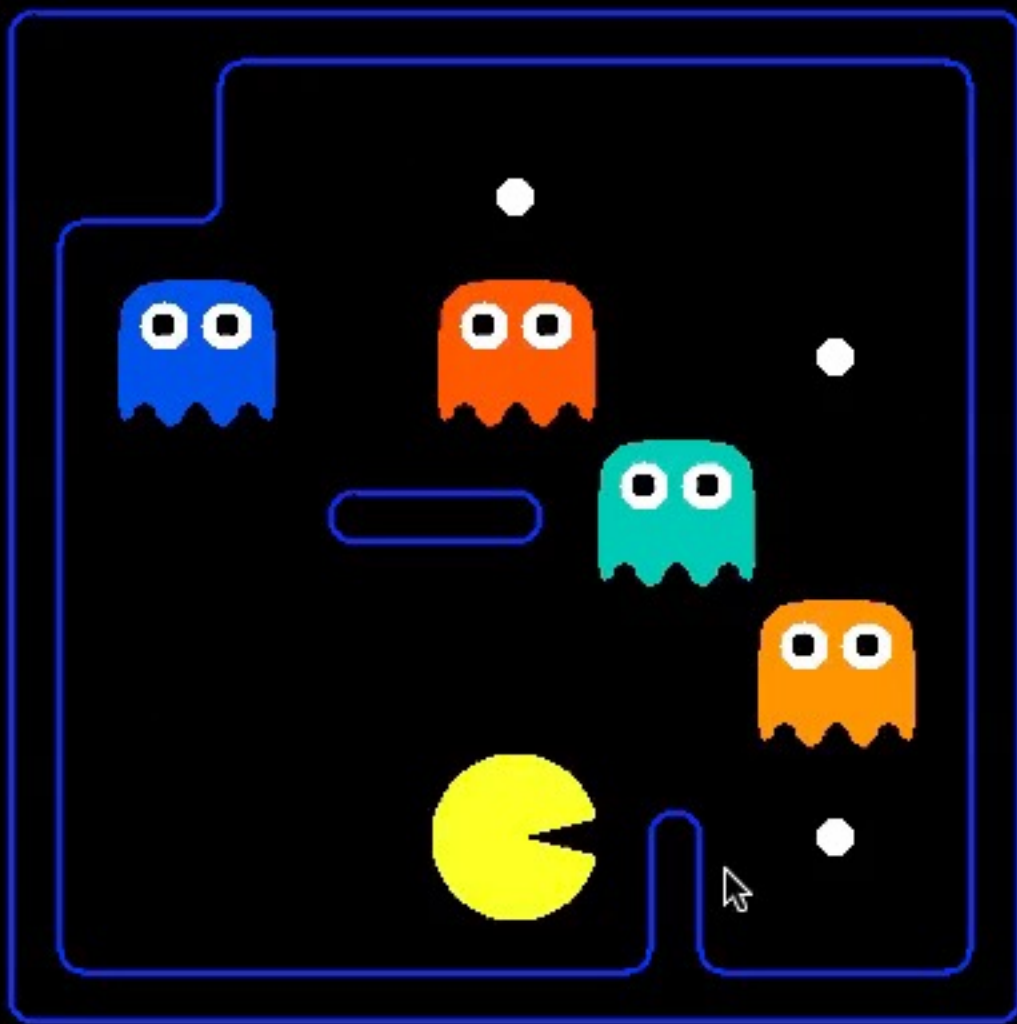
# How many sentences?

- Vast majority of KB occupied by  $O(NT)$  transition model sentences
  - Each about 10 lines of text
  - $N=200, T=400 \Rightarrow \sim 800,000$  lines of text, or 20,000 pages
- This is because propositional logic has limited expressive power
- Are we really going to write 20,000 pages of logic sentences???
- No, but your code will generate all those sentences!
- In first-order logic, we need  $O(1)$  transition model sentences
- (State-space search uses atomic states: how do we keep the transition model representation small???)

# Some reasoning tasks

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- **Localization** with a map and local sensing:
  - Given an initial KB, plus a sequence of percepts and actions, where am I?
- **Mapping** with a location sensor:
  - Given an initial KB, plus a sequence of percepts and actions, what is the map?
- **Simultaneous localization and mapping**:
  - Given ..., where am I and what is the map?
- **Planning**:
  - Given ..., what action sequence is guaranteed to reach the goal?
- **ALL OF THESE USE THE SAME KB AND THE SAME ALGORITHM!!**



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# Summary

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- One possible agent architecture: knowledge + inference
- Logics provide a formal way to encode knowledge
  - A logic is defined by: syntax, set of possible worlds, truth condition
- A simple KB for Pacman covers the initial state, sensor model, and transition model
- Logical inference computes entailment relations among sentences, enabling a wide range of tasks to be solved