4.6 Harrow-Hassidim-Lloyd algorithm

The Harrow-Hassidim-Lloyd (HHL) algorithm allows for the resolution of linear system problems on a quantum computer. To be precise, the problem to be solved is described as finding the N_b complex entries of \bf{x} that solve the following problem

$$
A\mathbf{x} = \mathbf{b},\tag{4.61}
$$

where *A* is an hermitian and non-singular $N_b \times N_b$ matrix and **b** is a N_b vector, both defined on \mathbb{C} . Classically, the solution is given by

$$
\mathbf{x} = A^{-1}\mathbf{b}.\tag{4.62}
$$

The question is then how one can implement this on a quantum computer.

First, let us assume that the entries of **b** are such that $||\mathbf{b}|| = 1$. Then, **b** can be stored in a n_b -qubit state $|b\rangle$, through the following mapping:

$$
\mathbf{b} = \begin{pmatrix} b_0 \\ \vdots \\ b_{N_b - 1} \end{pmatrix} \leftrightarrow b_0 |0\rangle + \dots + b_{N_b - 1} |N_b - 1\rangle = |b\rangle, \qquad (4.63)
$$

where $N_b = 2^{n_b}$. For example, this can be done via a unitary operation \hat{U}_b . Now, we define $|x\rangle = \hat{A}^{-1} |b\rangle$, where \hat{A} in the computational representation gives the classical matrix A. Notably, the state $|x\rangle$ needs to be normalised to be stored in a quantum register. Thus, one has

$$
|x\rangle = \frac{\hat{A}^{-1} |b\rangle}{||\hat{A}^{-1} |b\rangle||},\tag{4.64}
$$

where the normalisation problem can be tackled in a second moment.

Consider the spectral decomposition of *A*ˆ:

$$
\hat{A}|v_j\rangle = \lambda_j |v_j\rangle, \qquad (4.65)
$$

where λ_j and $|v_j\rangle$ are respectively the eigeinvalues and eigeinstates of \hat{A} . We also assume that the ordering of the eigeinvalues is such that

$$
0 < \lambda_0 \le \dots \le \lambda_{N_b - 1} < 1. \tag{4.66}
$$

In general this will not be the case, but one can remap the problem in order to fall within this case. We also assume that all the N_b eigeinvalues have an exact d -bit representation.

By applying what in Sec. $\boxed{4.5}$, we can query \hat{A} via an unitary operation $\hat{U} = e^{2\pi i \hat{A}}$ using QPE. For example, suppose $|b\rangle = |v_j\rangle$, then we have

$$
\hat{U}_{\text{QPE}}\left|0\right\rangle^{\otimes d}\left|v_{j}\right\rangle = \left|\lambda_{j}\right\rangle\left|v_{j}\right\rangle. \tag{4.67}
$$

In particular, the (not-normalised) solution of the linear system problem would be

$$
\hat{A}^{-1} |b\rangle = \hat{A}^{-1} |v_j\rangle = \frac{1}{\lambda_j} |v_j\rangle.
$$
\n(4.68)

More generally, one can decompose the state $|b\rangle$ on the basis of \hat{A} , i.e.

$$
|b\rangle = \sum_{j=0}^{2^{n_b}-1} \beta_j |v_j\rangle, \qquad (4.69)
$$

where β_j are a linear combination of b_j . Then the QPE procedure gives

52 4 Circuit model for quantum computation

$$
\hat{U}_{\text{QPE}}\left|0\right\rangle^{\otimes d}\left|b\right\rangle = \sum_{j} \beta_{j} \left|\lambda_{j}\right\rangle \left|v_{j}\right\rangle, \tag{4.70}
$$

and the solution of the problem is given by

$$
\hat{A}^{-1} |b\rangle = \sum_{j=0}^{2^{n_b}-1} \frac{\beta_j}{\lambda_j} |v_j\rangle.
$$
 (4.71)

The aim of the HHL algorithm is to generate the normalised version of the state in Eq. (4.71) from the general state $|b\rangle$ as shown in Eq. (4.69) .

The algorithm works with three registers. The first one is an ancillary register made of a single qubit, the second is also an ancillary register but made of *d* qubits, the third register is made of *n^b* qubits and will encode the solution of the problem. The HHL circuit is the following

$$
|0\rangle \longrightarrow \frac{| \Psi_1 \rangle | \Psi_4 \rangle | \Psi_5 \rangle | \Psi_6 \rangle | \Psi_9 \rangle
$$

\n
$$
|0\rangle^{\otimes d} \longrightarrow \frac{1}{| \Psi_1 | \Psi_2 |} \frac{1}{| \Psi_3 |} \frac{1}{| \Psi_4 |} \frac{1}{| \Psi_5 |} \frac{1}{| \Psi_6 |} \frac{1}{| \Psi_7 |} \frac{1}{| \Psi_8 |} \frac{1}{| \Psi_9 |} \frac{1}{| \Psi_9 |} \frac{1}{| \Psi_1 |} \frac{1}{| \Psi_2 |} \frac{1}{| \Psi_5 |} \frac{1}{| \Psi_7 |} \frac{1}{| \Psi_8 |} \frac{1}{| \Psi_9 |} \frac{
$$

The algorithm works as the following. Initially, all the qubits are prepared in $|0\rangle$:

$$
|\Psi_0\rangle = |0\rangle |0\rangle^{\otimes d} |0\rangle^{\otimes n_b},\qquad(4.73)
$$

then the information about b is encoded in the last register:

$$
|\Psi_1\rangle = \hat{\mathbb{1}} \otimes \hat{\mathbb{1}}^{\otimes d} \otimes \hat{U}_{\mathbf{b}} |\Psi_0\rangle = |0\rangle |0\rangle^{\otimes d} |b\rangle.
$$
 (4.74)

We apply the QPE procedure, which is here broke down in the corresponding three steps. The first is the application of the Hadamard gate:

$$
|\Psi_2\rangle = \hat{\mathbb{1}} \otimes \hat{H}^{\otimes d} \otimes \hat{\mathbb{1}} |\Psi_1\rangle = |0\rangle \frac{1}{2^{d/2}} (|0\rangle + |1\rangle)^{\otimes d} |b\rangle.
$$
 (4.75)

This is followed by the controlled unitary \hat{U}^{j} :

$$
|\Psi_3\rangle = \hat{\mathbb{1}} \otimes C(U^j) |\Psi_2\rangle = |0\rangle \frac{1}{2^{d/2}} \sum_{k=0}^{2^{d}-1} e^{2\pi ik\varphi} |k\rangle |b\rangle , \qquad (4.76)
$$

where $\hat{U} |b\rangle = e^{2\pi i \varphi} |b\rangle$ with $\varphi \in [0,1]$. Finally, we apply the inverse Fourier transform to the second register

$$
\begin{split} |\Psi_4\rangle &= \hat{\mathbb{1}} \otimes \hat{F}^\dagger \otimes \hat{\mathbb{1}}^{\otimes n_b} | \Psi_3 \rangle \,, \\ &= |0\rangle \frac{1}{2^{d/2}} \sum_{k=0}^{2^{d}-1} e^{2\pi ik\varphi} \hat{F}^\dagger | k \rangle | b \rangle \,, \\ &= |0\rangle \frac{1}{2^d} \sum_{k=0}^{2^d-1} e^{2\pi ik\varphi} \sum_{y=0}^{2^d-1} e^{-2\pi i y k/2^d} | y \rangle | b \rangle \,. \end{split} \tag{4.77}
$$

4.6 Harrow-Hassidim-Lloyd algorithm 53

However, one has that

$$
\sum_{k=0}^{2^d-1} e^{2\pi i k (\varphi - y/2^d)} = \begin{cases} \sum_{k=0}^{2^d-1} e^0 = 2^d, & \text{if } \varphi = y/2^d, \\ 0, & \text{if } \varphi \neq y/2^d, \end{cases}
$$
(4.78)

meaning that the *k* sum selects the value of $y = \varphi 2^d$. Thus,

$$
|\Psi_4\rangle = |0\rangle |\varphi_2^d\rangle |b\rangle. \tag{4.79}
$$

In general, $|b\rangle$ is in a superposition of $|v_j\rangle$, then

$$
\hat{U} |v_j\rangle = e^{2\pi i \hat{A}} |v_j\rangle = e^{2\pi i \lambda_j} |v_j\rangle.
$$
\n(4.80)

Then, the entire QPE gate maps

$$
|\Psi_1\rangle = |0\rangle |0\rangle^{\otimes d} \sum_{j=0}^{2^{n_b}-1} \beta_j |v_j\rangle \xrightarrow{U_{\text{QPE}}} |\Psi_4\rangle = |0\rangle \sum_{j=0}^{2^{n_b}-1} \beta_j | \lambda_j 2^d \rangle |v_j\rangle. \tag{4.81}
$$

We apply a controlled rotation on the first register, such that

$$
|\Psi_5\rangle = C(R) \otimes \hat{\mathbb{1}}^{\otimes n_b} |\Psi_4\rangle = \sum_{j=0}^{2^{n_b}-1} \beta_j \left(\sqrt{1 - \frac{C^2}{\lambda_j^2}} |0\rangle + \frac{C}{\lambda_j} |1\rangle \right) |\lambda_j 2^d\rangle |v_j\rangle, \qquad (4.82)
$$

where $C \in \mathbb{R}$ is an arbitrary constant. At this point we perform the measurement of the first register. If the outcome is $+1$ and the state collapses in $|0\rangle$ then we discard the run; if the outcome is -1 with the state collapsed in $|1\rangle$ then we retain the run. To increase the probabilities of having the outcome -1 , we make C as large as possible. After the collapse of the first register in $|1\rangle$, the state of the second and third register is

$$
|\Psi_6\rangle = \frac{1}{\left(\sum_{j=0}^{2^{n_b}-1} |\beta_j/\lambda_j|^2\right)^{1/2}} \sum_{j=0}^{2^{n_b}-1} \frac{\beta_j}{\lambda_j} |\lambda_j 2^d\rangle |v_j\rangle, \qquad (4.83)
$$

where we exploited that $C \in \mathbb{R}$. Now, we apply the inverse QPE, which has also three steps. The first is the application of the QFT:

$$
\begin{split}\n\ket{\Psi_7} &= \hat{F} \otimes \hat{\mathbb{1}}^{\otimes n_b} \ket{\Psi_6}, \\
&= \frac{1}{\left(\sum_{j=0}^{2^{n_b}-1} |\beta_j/\lambda_j|^2\right)^{1/2}} \sum_{j=0}^{2^{n_b}-1} \frac{\beta_j}{\lambda_j} \hat{F} \ket{\lambda_j 2^d} \ket{v_j}, \\
&= \frac{1}{\left(\sum_{j=0}^{2^{n_b}-1} |\beta_j/\lambda_j|^2\right)^{1/2}} \sum_{j=0}^{2^{n_b}-1} \frac{\beta_j}{\lambda_j} \frac{1}{2^{d/2}} \sum_{y=0}^{2^{d}-1} e^{2\pi i y (\lambda_j 2^d)/2^d} \ket{y} \ket{v_j}.\n\end{split} \tag{4.84}
$$

Then, we apply the controlled unitary $C(U^{-j})$, which gives

$$
|\Psi_8\rangle = C(U^{-j}) |\Psi_7\rangle, = \frac{1}{\left(\sum_{j=0}^{2^{n_b}-1} |\beta_j/\lambda_j|^2\right)^{1/2}} \sum_{j=0}^{2^{n_b}-1} \frac{\beta_j}{\lambda_j} \frac{1}{2^{d/2}} \sum_{y=0}^{2^{d}-1} e^{2\pi i y \lambda_j} |y\rangle e^{-2\pi i \lambda_j y} |v_j\rangle,
$$
(4.85)

where the two phases cancel and thus

54 4 Circuit model for quantum computation

$$
|\Psi_8\rangle = \frac{1}{2^{d/2}} \sum_{y=0}^{2^{d}-1} |y\rangle \sum_{j=0}^{2^{n_b}-1} \frac{\beta_j}{\lambda_j} \frac{1}{\left(\sum_{j=0}^{2^{n_b}-1} |\beta_j/\lambda_j|^2\right)^{1/2}} |v_j\rangle.
$$
 (4.86)

Finally, the application of Hadamard's gates on the second register gives

$$
|\Psi_9\rangle = \hat{H}^{\otimes d} \otimes \hat{\mathbb{1}}^{\otimes n_b} |\Psi_8\rangle = |0\rangle^{\otimes d} \sum_{j=0}^{2^{n_b}-1} \frac{\beta_j}{\lambda_j} \frac{1}{\left(\sum_{j=0}^{2^{n_b}-1} |\beta_j/\lambda_j|^2\right)^{1/2}} |v_j\rangle, \tag{4.87}
$$

where the third register is exactly in the form in Eq. (4.71) after the proper normalisation. Thus,

$$
|\Psi_9\rangle = |0\rangle^{\otimes d} |x\rangle \,,\tag{4.88}
$$

embeds the solution of the linear system $A\mathbf{x} = \mathbf{b}$.