4.6 Harrow-Hassidim-Lloyd algorithm

The Harrow-Hassidim-Lloyd (HHL) algorithm allows for the resolution of linear system problems on a quantum computer. To be precise, the problem to be solved is described as finding the N_b complex entries of **x** that solve the following problem

$$A\mathbf{x} = \mathbf{b},\tag{4.61}$$

where A is an hermitian and non-singular $N_b \times N_b$ matrix and **b** is a N_b vector, both defined on \mathbb{C} . Classically, the solution is given by

$$\mathbf{x} = A^{-1}\mathbf{b}.\tag{4.62}$$

The question is then how one can implement this on a quantum computer.

First, let us assume that the entries of **b** are such that $||\mathbf{b}|| = 1$. Then, **b** can be stored in a n_b -qubit state $|b\rangle$, through the following mapping:

$$\mathbf{b} = \begin{pmatrix} b_0 \\ \vdots \\ b_{N_b-1} \end{pmatrix} \leftrightarrow b_0 |0\rangle + \dots + b_{N_b-1} |N_b - 1\rangle = |b\rangle, \qquad (4.63)$$

where $N_b = 2^{n_b}$. For example, this can be done via a unitary operation $\hat{U}_{\mathbf{b}}$. Now, we define $|x\rangle = \hat{A}^{-1} |b\rangle$, where \hat{A} in the computational representation gives the classical matrix A. Notably, the state $|x\rangle$ needs to be normalised to be stored in a quantum register. Thus, one has

$$|x\rangle = \frac{\hat{A}^{-1}|b\rangle}{||\hat{A}^{-1}|b\rangle||},\tag{4.64}$$

where the normalisation problem can be tackled in a second moment.

Consider the spectral decomposition of \hat{A} :

$$\hat{A} |v_j\rangle = \lambda_j |v_j\rangle, \qquad (4.65)$$

where λ_j and $|v_j\rangle$ are respectively the eigenvalues and eigenstates of \hat{A} . We also assume that the ordering of the eigenvalues is such that

$$0 < \lambda_0 \le \dots \le \lambda_{N_b - 1} < 1. \tag{4.66}$$

In general this will not be the case, but one can remap the problem in order to fall within this case. We also assume that all the N_b eigenvalues have an exact *d*-bit representation.

By applying what in Sec. 4.5, we can query \hat{A} via an unitary operation $\hat{U} = e^{2\pi i \hat{A}}$ using QPE. For example, suppose $|b\rangle = |v_j\rangle$, then we have

$$\hat{U}_{\text{QPE}} \left| 0 \right\rangle^{\otimes d} \left| v_j \right\rangle = \left| \lambda_j \right\rangle \left| v_j \right\rangle. \tag{4.67}$$

In particular, the (not-normalised) solution of the linear system problem would be

$$\hat{A}^{-1} |b\rangle = \hat{A}^{-1} |v_j\rangle = \frac{1}{\lambda_j} |v_j\rangle.$$
(4.68)

More generally, one can decompose the state $|b\rangle$ on the basis of \hat{A} , i.e.

$$|b\rangle = \sum_{j=0}^{2^{n_b}-1} \beta_j |v_j\rangle, \qquad (4.69)$$

where β_j are a linear combination of b_j . Then the QPE procedure gives

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$$\hat{U}_{\text{QPE}} \left| 0 \right\rangle^{\otimes d} \left| b \right\rangle = \sum_{j} \beta_{j} \left| \lambda_{j} \right\rangle \left| v_{j} \right\rangle, \tag{4.70}$$

and the solution of the problem is given by

$$\hat{A}^{-1} |b\rangle = \sum_{j=0}^{2^{n_b}-1} \frac{\beta_j}{\lambda_j} |v_j\rangle.$$
(4.71)

The aim of the HHL algorithm is to generate the normalised version of the state in Eq. (4.71) from the general state $|b\rangle$ as shown in Eq. (4.69).

The algorithm works with three registers. The first one is an ancillary register made of a single qubit, the second is also an ancillary register but made of d qubits, the third register is made of n_b qubits and will encode the solution of the problem. The HHL circuit is the following

The algorithm works as the following. Initially, all the qubits are prepared in $|0\rangle$:

$$|\Psi_0\rangle = |0\rangle |0\rangle^{\otimes d} |0\rangle^{\otimes n_b}, \qquad (4.73)$$

then the information about \mathbf{b} is encoded in the last register:

$$|\Psi_1\rangle = \hat{\mathbb{1}} \otimes \hat{\mathbb{1}}^{\otimes d} \otimes \hat{U}_{\mathbf{b}} |\Psi_0\rangle = |0\rangle |0\rangle^{\otimes d} |b\rangle.$$
(4.74)

We apply the QPE procedure, which is here broke down in the corresponding three steps. The first is the application of the Hadamard gate:

$$|\Psi_2\rangle = \hat{\mathbb{1}} \otimes \hat{H}^{\otimes d} \otimes \hat{\mathbb{1}} |\Psi_1\rangle = |0\rangle \, \frac{1}{2^{d/2}} (|0\rangle + |1\rangle)^{\otimes d} |b\rangle \,. \tag{4.75}$$

This is followed by the controlled unitary \hat{U}^j :

$$|\Psi_3\rangle = \hat{\mathbb{1}} \otimes C(U^j) |\Psi_2\rangle = |0\rangle \frac{1}{2^{d/2}} \sum_{k=0}^{2^d-1} e^{2\pi i k\varphi} |k\rangle |b\rangle, \qquad (4.76)$$

where $\hat{U} |b\rangle = e^{2\pi i \varphi} |b\rangle$ with $\varphi \in [0, 1[$. Finally, we apply the inverse Fourier transform to the second register

$$\begin{aligned} |\Psi_{4}\rangle &= \mathbb{1} \otimes F^{\dagger} \otimes \mathbb{1}^{\otimes n_{b}} |\Psi_{3}\rangle, \\ &= |0\rangle \frac{1}{2^{d/2}} \sum_{k=0}^{2^{d}-1} e^{2\pi i k\varphi} \hat{F}^{\dagger} |k\rangle |b\rangle, \\ &= |0\rangle \frac{1}{2^{d}} \sum_{k=0}^{2^{d}-1} e^{2\pi i k\varphi} \sum_{y=0}^{2^{d}-1} e^{-2\pi i y k/2^{d}} |y\rangle |b\rangle. \end{aligned}$$

$$(4.77)$$

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However, one has that

$$\sum_{k=0}^{2^{d}-1} e^{2\pi i k (\varphi - y/2^{d})} = \begin{cases} \sum_{k=0}^{2^{d}-1} e^{0} = 2^{d}, & \text{if } \varphi = y/2^{d}, \\ 0, & \text{if } \varphi \neq y/2^{d}, \end{cases}$$
(4.78)

meaning that the k sum selects the value of $y = \varphi 2^d$. Thus,

$$|\Psi_4\rangle = |0\rangle |\varphi 2^d\rangle |b\rangle. \tag{4.79}$$

In general, $|b\rangle$ is in a superposition of $|v_j\rangle$, then

$$\hat{U}|v_j\rangle = e^{2\pi i \hat{A}}|v_j\rangle = e^{2\pi i \lambda_j}|v_j\rangle.$$
(4.80)

Then, the entire QPE gate maps

$$\left|\Psi_{1}\right\rangle = \left|0\right\rangle\left|0\right\rangle^{\otimes d} \sum_{j=0}^{2^{n_{b}}-1} \beta_{j} \left|v_{j}\right\rangle \xrightarrow{U_{\text{QPE}}} \left|\Psi_{4}\right\rangle = \left|0\right\rangle \sum_{j=0}^{2^{n_{b}}-1} \beta_{j} \left|\lambda_{j} 2^{d}\right\rangle\left|v_{j}\right\rangle.$$

$$(4.81)$$

We apply a controlled rotation on the first register, such that

$$|\Psi_5\rangle = C(R) \otimes \hat{\mathbb{1}}^{\otimes n_b} |\Psi_4\rangle = \sum_{j=0}^{2^{n_b}-1} \beta_j \left(\sqrt{1 - \frac{C^2}{\lambda_j^2}} |0\rangle + \frac{C}{\lambda_j} |1\rangle \right) |\lambda_j 2^d\rangle |v_j\rangle, \qquad (4.82)$$

where $C \in \mathbb{R}$ is an arbitrary constant. At this point we perform the measurement of the first register. If the outcome is +1 and the state collapses in $|0\rangle$ then we discard the run; if the outcome is -1 with the state collapsed in $|1\rangle$ then we retain the run. To increase the probabilities of having the outcome -1, we make C as large as possible. After the collapse of the first register in $|1\rangle$, the state of the second and third register is

$$|\Psi_{6}\rangle = \frac{1}{\left(\sum_{j=0}^{2^{n_{b}}-1} |\beta_{j}/\lambda_{j}|^{2}\right)^{1/2}} \sum_{j=0}^{2^{n_{b}}-1} \frac{\beta_{j}}{\lambda_{j}} |\lambda_{j}2^{d}\rangle |v_{j}\rangle, \qquad (4.83)$$

where we exploited that $C \in \mathbb{R}$. Now, we apply the inverse QPE, which has also three steps. The first is the application of the QFT:

$$\begin{aligned} |\Psi_{7}\rangle &= \hat{F} \otimes \hat{1}^{\otimes n_{b}} |\Psi_{6}\rangle ,\\ &= \frac{1}{\left(\sum_{j=0}^{2^{n_{b}-1}} |\beta_{j}/\lambda_{j}|^{2}\right)^{1/2}} \sum_{j=0}^{2^{n_{b}-1}} \frac{\beta_{j}}{\lambda_{j}} \hat{F} |\lambda_{j}2^{d}\rangle |v_{j}\rangle ,\\ &= \frac{1}{\left(\sum_{j=0}^{2^{n_{b}-1}} |\beta_{j}/\lambda_{j}|^{2}\right)^{1/2}} \sum_{j=0}^{2^{n_{b}-1}} \frac{\beta_{j}}{\lambda_{j}} \frac{1}{2^{d/2}}}{\sum_{y=0}^{2^{d}-1}} e^{2\pi i y (\lambda_{j}2^{d})/2^{d}}} |y\rangle |v_{j}\rangle . \end{aligned}$$
(4.84)

Then, we apply the controlled unitary $C(U^{-j})$, which gives

$$\begin{aligned} |\Psi_8\rangle &= C(U^{-j}) |\Psi_7\rangle, \\ &= \frac{1}{\left(\sum_{j=0}^{2^{n_b}-1} |\beta_j/\lambda_j|^2\right)^{1/2}} \sum_{j=0}^{2^{n_b}-1} \sum_{j=0}^{\beta_j} \frac{1}{\lambda_j} \frac{2^{d}-1}{2^{d/2}} \sum_{y=0}^{2^{d}-1} e^{2\pi i y \lambda_j} |y\rangle e^{-2\pi i \lambda_j y} |v_j\rangle, \end{aligned}$$
(4.85)

where the two phases cancel and thus

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$$|\Psi_8\rangle = \frac{1}{2^{d/2}} \sum_{y=0}^{2^d-1} |y\rangle \sum_{j=0}^{2^{n_b}-1} \frac{\beta_j}{\lambda_j} \frac{1}{\left(\sum_{j=0}^{2^{n_b}-1} |\beta_j/\lambda_j|^2\right)^{1/2}} |v_j\rangle.$$
(4.86)

Finally, the application of Hadamard's gates on the second register gives

$$|\Psi_{9}\rangle = \hat{H}^{\otimes d} \otimes \hat{1}^{\otimes n_{b}} |\Psi_{8}\rangle = |0\rangle^{\otimes d} \sum_{j=0}^{2^{n_{b}}-1} \frac{\beta_{j}}{\lambda_{j}} \frac{1}{\left(\sum_{j=0}^{2^{n_{b}}-1} |\beta_{j}/\lambda_{j}|^{2}\right)^{1/2}} |v_{j}\rangle,$$
(4.87)

where the third register is exactly in the form in Eq. (4.71) after the proper normalisation. Thus,

$$\left|\Psi_{9}\right\rangle = \left|0\right\rangle^{\otimes d}\left|x\right\rangle,\tag{4.88}$$

embeds the solution of the linear system $A\mathbf{x} = \mathbf{b}$.

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