

Classical predicate logic

A brief reminder

- Propositional logic is not enough to account for all valid arguments in natural language.
- For example:

*All humans are mortal,  
Socrates is human,  
therefore, Socrates is mortal.*

Cannot be accounted for in propositional logic (Try!)

- To cover these arguments a richer language is needed.
- We need a language able to analyse the internal structure of sentences (beyond just P or Q) and capturing quantifiers such as “all” and “some”.

- This is done by predicate logic.

In what follows a quick review of predicate logic is given.

The formal language of classical predicate logic

- The classical predicate language  $L_1$  is given by extending the classical propositional language  $L$ .
- We extend: the vocabulary, the syntax, the semantics of the propositional language  $L$ .

# Vocabulary (extension)

- We have two kinds of new expressions:
- **Individual constants:**  $p, q, r, s, \dots$  (like names of individuals)
- **Individual variables:**  $x, y, z, \dots$  (like names of unspecified individuals)
- **N-place predicates:**  $P(\dots), Q(\dots), \dots$  (like expressions for properties and relations)



- **Quantifiers:**

$\forall$  (The universal quantifier: *For all*)

$\exists$  (The existential quantifier: *For some/There exists at least one*)

# Syntax of L1 (Extension)

- The syntax of  $L_1$  is obtained by adding the following clauses to the propositional syntax.

**1.**

If  $A( )$  is a predicate symbols with  $n$ -places, then writing  $n$  individual constants or variables, separated by commas, in  $A( )$ , what is obtained a WFF.

**2.**

If  $A$  is a WFF with no occurrences of  $\forall v$  or  $\exists v$  for the variable  $v$ ,

Then  $\forall v A$  and  $\exists v A$  are WFF.

- The syntactic notions are defined as in L.
- An additional notion is that of **free variable** and **bound variable**.

- A variable  $v$  is free in wff if it is not in the range of a quantifier.

It is bounded otherwise.

- A wff is **closed** if it no free variables occur in it. Otherwise the formula is **open**.

(Note: these definition are intended to be intuitive reminder, but they are not fully rigorous. For full definition I refer to standard textbooks introducing predicate logic).

- Closed wffs are also called **sentences**.
- We only consider sentences (closed formulas) from now on.



# Semantics

- The semantics for the predicate language is much more complex than the semantics for the propositional language.

- In particular, truth tables do not extend to the new expressions (predicates, individual constants, variables, quantifiers).

- To fix these a more complex semantics, based on the notions like that of domain, interpretation, assignment, satisfaction, model, ... must be introduced.

- Here, however, we just give an intuitive semantic reading for the new expressions, and then turn to semantic trees.

- Notice that, unlike truth tables, refutation trees can be easily extended to the language of predicate logic.

→ And also to modal predicate logic.

We thus follow that route.

- Intuitively, the domain is a set of objects. The set of objects we intend to speak about.
- Interpretations give meaning relatively to that set of objects.

- **Individual constants** are given a meaning by associating to every constant an object of the domain.



- **Individual variables** are assigned a “temporary” and variable objects by various assignments.
  - A complete treatment would be needed to give rigorous semantic clauses for quantifiers.

- A **predicate symbol** with  $n$  individual constants is read as meaning that a certain relation hold among those  $n$  individuals.

For example  $P(t,r)$  means that  $t$  and  $r$  are in the relation  $P$ .

(e.g. if  $P$  stands for “loves”,  $t$  loves  $r$ ).

- For simplicity, we limit our treatment to unary predicate symbols ascribing a property to certain objects.
- For example,  $Q(r)$  means that  $r$  is  $Q$ .

- $\forall xQ(x)$  means that all objects  $x$  are  $Q$
- $\exists yR(y)$  means that at least one object is  $R$

- In complete treatment this should be phrased in terms of assignments of values to individual variables.
- The details are not obvious and we skip them here.

- Note that:

*All humans are mortal* can be translated as:

$$\forall y (H(y) \rightarrow M(y)) \quad \forall \quad \exists$$

Namely: for all  $y$ , if  $y$  is human, then  $y$  is mortal.

With a **conditional!**

*Some humans are mortal* can be translated as:

$$\exists y(H(y) \& M(y))$$

Namely: there is at least one  $y$ , such that  $y$  is human and  $y$  is mortal.

With a **conjunction**!

- For other notable translations, such as *no human is mortal*, see an introduction to predicate logic.





Semantics

Appendix

Note:

What follows is an additional appendix not required for the course on modal logic, and just put as material for the interested reader.

- A **little** more rigorously, the semantics for the predicate language could be given as follows.

- For predicative logic, we need a more complex structure (also called a model) consisting of:
  1. A domain  $D$  and
  2. An interpretation  $I$ .

- D is a set of element.
- It is intended to provide the things we talk about.
- I is the semantic interpretation that gives meaning to the expression of L1.  
I interprets L1 in D.

- In particular, for the basic vocabulary:

I gives an element to each individual constant.

I gives a set of elements to unary predicates (the set of things that have the intended property)

- Then we can define **truth in the model  $M = (D, I)$**  as follows. The definition is inductive:

*Base of induction:*

An atomic sentence, like  $P(a)$ , is true in  $M$  if the element given to  $a$  by  $I$  is in the set given to  $P()$  by  $I$ .



Inductive steps:

- Propositional combinations of formulas follows truth tables.

(so  $A \& B$  is true in  $M$  iff  $A$  is true in  $M$  and  $B$  is true in  $M$ )

- For the quantifiers, we have that

$\forall vA$  is true in  $M$ , if everything in  $M$  has the property  $A$ .

$\exists vA$  is true in  $M$ , if something in  $M$  has the property  $A$ .

- With the notion of truth in  $M$ , we can define logical validity as usual.
- An argument is logically validity when:  
if the premises are true in every model, the conclusion is true in every model.

- Notice that the above is not the fully rigorous treatment, which require the notion of satisfaction and assignment of values to variables.
- For such a treatment, refer, again, to a textbook in predicate logic.

# Semantics for modal predicate logic



- The semantics works like in predicate logic, interpreting and defining truth in a model ( $M = D, I$ ).
- Additionally however,  $M$  specifies a set  $W$ , the set of “possible worlds” of  $M$ ,  
one of which is designated its “actual world”,  
and each world  $w$  in  $W$  is assigned its own domain of quantification,  
 **$d(w)$  included in  $D$**  (intuitively, the set of individuals that exist in  $w$ ).

- For each basic expression (constant and predicate) we have an **extension**, relative to each world,  
and an **intension**. (the function that assigns the extension).



- The definition of truth is the same of predicate logic but relativized to possible worlds.

So we have **Truth in M, relatively to a world w.**

- For example, an atomic sentence like  $P(a)$  is true in  $M$ , in world  $w$  iff  
the extension of  $a$  in  $w$  belongs to the extension of  $P()$  in  $w$ .
- The clauses for connective remains the same.

- For quantifiers, we need to specify that quantifiers are limited to a single world, and they do not run on the entire domain (given by the union of all elements in all possible worlds).

$\forall vA$  is true in  $M, w$ , if everything in  $w$  has the property  $A$ .

$\exists vA$  is true in  $M, w$  if something in  $w$  has the property  $A$ .

- The clauses of modal operators also remains the same, just relativized to possible worlds.
  - $\Box P$  is true in  $M, w$  if and only if  $p$  is true in all possible worlds.
  - $\Diamond P$  is true in  $M, w$  if and only if  $p$  is true in some possible worlds.

- We can then define logical validity, as usual just relativized to possible worlds.

*De dicto* and *de re*

- An application of the possible world analysis concerns the venerable distinction between *de re* and *de dicto* modality.
- Among the strongest modal intuitions is that the possession of a property has a modal character — that things exemplify some properties necessarily, or essentially, and others only accidentally.



- We can say that an individual  $a$  has a property  $P$  essentially if  $a$  has  $P$  in every world in which it exists.

Algol is a dog essentially:  $\Box G(a)$

Algol is a pet accidentally:  $Ta \ \& \ \Diamond \neg T(a)$

Sentences like these, in which properties are ascribed to a specific individual in a modal context, signaled formally by the occurrence **of a name or the free occurrence of a variable** in the scope of a modal operator, are de re.

- Often, *de re* statements are identified, more strictly, with those in which there is quantification in the scope of modal operator.

So, there is a free variable in the scope.

- These are easily obtained by existential generalizations, even if individual constants are involved.

So, this sense includes the former.

- Necessarily, all dogs are mammals:

$$\Box \forall x(D(x) \rightarrow M(x))$$

$$\Box (P \vee Q)$$

Where in the scope of the modal operators only bound variables, at most, occur are *de dicto*.

- Note that propositional modal logic only expresses *de dicto* modalities.

- Sometimes, de re statement can be derived from *de dicto*, by logical steps, in modal predicate logic.
- This is a problem for those, like Quine, who rejects de re statements.

Indeed, Quine only accepts modal propositional logic at most.

- Modality *de re* involves in addition a commitment to the meaningfulness of ***transworld identity***.

The thesis that, necessarily, individuals exist and exemplify (often very different) properties in many different possible worlds.

- More specifically, basic possible world semantics:
  - (i) permits world domains to overlap
  - (ii) assigns intensions to predicates, thereby, in effect, relativizing predicate extensions to worlds
- Both features are absent in Lewis concretism.