



Binary Search Trees

Chapter 12 of Cormen's book

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Algorithmic Design

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Binary Search

Binary search is an efficient algorithm using a divide-and-conquer strategy. Its running time is $O(\log n)$.

Algorithm 3 Binary Search

- 1: **INPUT:** A sorted sequence $s = s[1]s[2] \dots s[n]$ of items from a set X and an item $x \in X$.
 - 2: **OUTPUT:** An index $i \in [1, n]$ such that $s[i] = x$; or **FAIL** if no such index exists.
 - 3: $\text{start} \leftarrow 1, \text{end} \leftarrow n$;
 - 4: **while** $\text{start} \leq \text{end}$ **do**
 - 5: $\text{mid} \leftarrow \lfloor (\text{start} + \text{end}) / 2 \rfloor$;
 - 6: **if** $s[\text{mid}] = x$ **then**
 - 7: **return** : mid ;
 - 8: **else if** $s[\text{mid}] < x$ **then**
 - 9: $\text{start} \leftarrow \text{mid} + 1$;
 - 10: **else if** $s[\text{mid}] > x$ **then**
 - 11: $\text{end} \leftarrow \text{mid} - 1$;
 - 12: **return** : **FAIL**;
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Binary Search Trees

BSTs have the following property:

For every node x in the tree, for every node y in the **left** subtree of x , then $key(x) \leq key(y)$; for every node z in the **right** subtree of x , $key(z) \geq key(x)$.

ITERATIVE-TREE-SEARCH(x, k)

```
1  while  $x \neq \text{NIL}$  and  $k \neq x.key$ 
2      if  $k < x.key$ 
3           $x = x.left$ 
4      else  $x = x.right$ 
5  return  $x$ 
```

Binary Search Trees

TREE-INSERT(T, z)

```
1   $y = \text{NIL}$ 
2   $x = T.\text{root}$ 
3  while  $x \neq \text{NIL}$ 
4       $y = x$ 
5      if  $z.\text{key} < x.\text{key}$ 
6           $x = x.\text{left}$ 
7      else  $x = x.\text{right}$ 
8   $z.p = y$ 
9  if  $y == \text{NIL}$ 
10      $T.\text{root} = z$            // tree  $T$  was empty
11  elseif  $z.\text{key} < y.\text{key}$ 
12      $y.\text{left} = z$ 
13  else  $y.\text{right} = z$ 
```

Binary Search Trees

INORDER-TREE-WALK(x)

1 **if** $x \neq \text{NIL}$

2 INORDER-TREE-WALK($x.\textit{left}$)

3 print $x.\textit{key}$

4 INORDER-TREE-WALK($x.\textit{right}$)

Binary Search Trees

TREE-MINIMUM(x)

```
1  while  $x.left \neq \text{NIL}$ 
2       $x = x.left$ 
3  return  $x$ 
```

TREE-MAXIMUM(x)

```
1  while  $x.right \neq \text{NIL}$ 
2       $x = x.right$ 
3  return  $x$ 
```

Binary Search Trees

TREE-SUCCESSOR(x)

```
1  if  $x.right \neq \text{NIL}$ 
2      return TREE-MINIMUM( $x.right$ )
3   $y = x.p$ 
4  while  $y \neq \text{NIL}$  and  $x == y.right$ 
5       $x = y$ 
6       $y = y.p$ 
7  return  $y$ 
```

Binary Search Trees

TRANSPLANT(T, u, v)

```
1  if  $u.p == \text{NIL}$ 
2       $T.root = v$ 
3  elseif  $u == u.p.left$ 
4       $u.p.left = v$ 
5  else  $u.p.right = v$ 
6  if  $v \neq \text{NIL}$ 
7       $v.p = u.p$ 
```


Binary Search Trees

TREE-DELETE(T, z)

- 1 **if** $z.left == \text{NIL}$
- 2 TRANSPLANT($T, z, z.right$)
- 3 **elseif** $z.right == \text{NIL}$
- 4 TRANSPLANT($T, z, z.left$)
- 5 **else** $y = \text{TREE-TREE-SUCCESSOR}(z)$
- 6 **if** $y.p \neq z$
- 7 TRANSPLANT($T, y, y.right$)
- 8 $y.right = z.right$
- 9 $y.right.p = y$
- 10 TRANSPLANT(T, z, y)
- 11 $y.left = z.left$
- 12 $y.left.p = y$

